47853 Packing and Covering: Lecture 5

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January 29, 2019

5 Perfect graphs

Let G = (V, E) be a simple graph. Recall that G is perfect if, for every induced subgraph G' of G, $\chi(G') = \omega(G')$. (Notice that G' may be G.) In words, a simple graph is perfect if in each induced subgraph, the maximum cardinality of a clique is equal to the chromatic number. Last time we proved the following:

Corollary 5.4. The following graphs are perfect:

- (1) bipartite graphs, and their complements,
- (2) line graphs of bipartite graphs, and their complements,
- (3) comparability graphs, and their complements.

The obvious question this corollary leads to is, does complementation preserve perfection? Claude Berge asked the same question in 1961 [1]. Today we will see that the answer is surprisingly yes!

5.1 The max-max inequality and the weak perfect graph theorem

The proof we present of the following result is due to Gasparian 1996 [4]:

Theorem 5.5 (Lovász 1972 [5]). Let G be a simple graph. The following statements are equivalent:

(i) G is perfect,

(ii) $\omega(H) \cdot \alpha(H) \ge |V(H)|$ for every induced subgraph H.

Proof. (i) \Rightarrow (ii): Let *H* be an induced subgraph. By definition, $\chi(H) = \omega(H)$, that is, V(H) can be covered by $\omega(H)$ stable sets. Since each stable set has cardinality at most $\alpha(H)$, it follows that

$$|V(H)| \le \omega(H) \cdot \alpha(H).$$

(ii) \Rightarrow (i): Suppose for a contradiction that G is not perfect. Let H be an induced subgraph of G that is not perfect, but every proper induced subgraph of H is perfect. Let $\omega := \omega(H)$, $\alpha := \alpha(H)$ and n := |V(H)|. Note that n > 1. Clearly,

 $\omega \ge \omega(H \setminus S) \ge \omega - 1$ for every nonempty stable set $S \subseteq V(H)$;

since $H \setminus S$ is perfect and H is not, it follows that

 $\omega(H \setminus S) = \omega$ for every nonempty stable set $S \subseteq V(H)$.

Let S_0 be a maximum stable set of H. Then for every vertex $v \in S_0$, $H \setminus v$ is perfect, so its vertices can be partitioned into $\omega(H \setminus v) = \omega$ nonempty stable sets. As S_0 has α vertices, we get $\alpha\omega$ stable sets $S_1, \ldots, S_{\alpha\omega}$.

Claim. Every maximum clique of H intersects all but one of $S_0, S_1, \ldots, S_{\alpha\omega}$ exactly once.

Proof of Claim. Let C be a maximum clique of H. Clearly C intersects each one of $S_0, S_1, \ldots, S_{\alpha\omega}$ at most once. For a vertex $v \in S_0$, if

- $v \in C$: then C intersects all but one stable set in every partition of $V(H \setminus v)$ into ω stable sets,
- $v \notin C$: then C intersects all stable sets in every partition of $V(H \setminus v)$ into ω stable sets.

This observation immediately implies the claim.

For each $i \in \{0, 1, ..., \alpha\omega\}$, let C_i be a maximum clique of $H \setminus S_i$; notice that $|C_i| = \omega$. Let A be the 0-1 matrix whose columns are labeled by V(H), and whose rows are the incidence vectors of $S_0, S_1, ..., S_{\alpha\omega}$. Let B be the 0-1 matrix whose columns are labeled by V(H), and whose rows are the incidence vectors of $C_0, C_1, ..., C_{\alpha\omega}$. It then follows from the claim above that $AB^{\top} = J - I$, where J is the all-ones matrix and I the identity matrix of appropriate dimensions. Since J - I is a nonsingular $(\alpha\omega + 1) \times (\alpha\omega + 1)$ matrix, it follows that both A and B have full row rank, implying in turn that

$$|V(H)| = n \ge \alpha \omega + 1 = \alpha(H) \cdot \omega(H) + 1 > |V(H)|,$$

a contradiction.

As a consequence, we get the weak perfect graph theorem:

Theorem 5.6 (Lovász 1972 [6]). If a graph is perfect, then so is its complement.

Proof. Suppose that G is perfect. Then by Theorem 5.5, for every induced subgraph H of G,

$$\omega(H) \cdot \alpha(H) \ge |V(H)|$$

implying in turn that for every induced subgraph \overline{H} of \overline{G} ,

$$\alpha(\overline{H}) \cdot \omega(\overline{H}) \ge |V(\overline{H})|,$$

so by Theorem 5.5, \overline{G} is perfect, as required.

 \Diamond

5.2 Odd holes and odd antiholes

We say that a simple graph is *minimally imperfect* if it is not perfect, but every proper induced subgraph is perfect. Equivalently, a simple graph G is minimally imperfect if $\chi(G) > \omega(G)$, but for every proper induced subgraph G', $\chi(G') = \omega(G')$. The latter implies that a minimally imperfect graph is always connected.

Remark 5.7. A graph is perfect if, and only if, it has no minimally imperfect induced subgraph.

Let *H* be an odd circuit with at least 5 vertices. Then $3 = \chi(H) > \omega(H) = 2$, so *G* is imperfect. Since every proper induced subgraph of *H* is bipartite, and therefore perfect, it follows that *H* is minimally imperfect. Notice that Theorem 5.6 equivalently states that,

Corollary 5.8. The complement of a minimally imperfect graph is also minimally imperfect.

Thus, the complement of an odd circuit with at least 5 vertices is also minimally imperfect. Let G be a simple graph. We say that G has an *odd hole* if it has as an induced subgraph an odd circuit with at least 5 vertices, and we say that G has an *odd antihole* if \overline{G} has an odd hole. It follows from the preceding remark that,

Remark 5.9. A perfect graph has no odd hole and no odd antihole.

In 1961, Claude Berge conjectured that the converse of this statement is also true [1]. In 2006, this conjecture was proved by Chudnovsky, Robertson, Seymour and Thomas, and their theorem is referred to as the *strong perfect graph theorem* [2]. We will see some of the milestones and highlights leading to the proof, as well as a sketch of the proof.

5.3 Star cutsets and antitwins

Let G = (V, E) be a simple graph. A *star cutset* is a nonempty $X \subseteq V$ such that

- $G \setminus X$ has more connected components than G, and
- a vertex of X is adjacent to all the other vertices in X.

Lemma 5.10 (Chvátal 1985 [3]). A minimally imperfect graph does not have a star cutset.

Proof. Let G = (V, E) be a minimally imperfect graph, and let $\omega := \omega(G)$. Then

 $\omega(G \setminus S) = \omega \quad \text{for every stable set } S \subseteq V.$

Suppose for a contradiction that G has a star cutset $X \subseteq V$. Then the vertices of $G \setminus X$ can be partitioned into nonempty parts V_1, V_2 such that G has no edge between V_1 and V_2 . Since every proper induced subgraph of G is perfect, for each $i \in [2]$, there is a vertex-coloring $f_i : X \cup V_i \to [\omega]$ of the induced subgraph $G[X \cup V_i]$. Since X is a star cutset, it has a vertex v that is adjacent to all other vertices of X. For $i \in [2]$, let $S_i := \{w \in X \cup V_i : f_i(w) = f_i(v)\}$. Clearly, each S_i is stable and $S_i \cap X = \{v\}$. Moreover, since there are no edges between V_1 and V_2 , it follows that $S := S_1 \cup S_2$ is also stable. In particular, $\omega(G \setminus S) = \omega$, so $G \setminus S$ has a clique C of cardinality ω . However, either $C \subseteq X \cup V_1$ or $C \subseteq X \cup V_2$, implying in turn that C is an ω -clique of some $G[X \cup V_i] \setminus S_i$, which has an $(\omega - 1)$ -vertex-coloring, a contradiction.

This lemma was a key milestone for what led to the proof of the strong perfect graph theorem. To demonstrate the power of this lemma, let us see some applications of it. Let G_1 be a perfect graph, and take a vertex $v \in V(G_1)$. To *duplicate* v is to introduce a new vertex \bar{v} , join it to all the neighbors of v, and then join it to \bar{v} . More generally, given another perfect graph G_2 over a disjoint vertex set, to *substitute* G_2 *for* v is to remove v, and join every vertex of G_2 to all the neighbors of v in $G_1 \setminus v$.

Theorem 5.11 (Lovász 1972 [6]). Let G_1, G_2 be perfect graphs over disjoint vertex sets. If G is obtained by substituting G_2 for a vertex v of G_1 , then G is perfect. In particular, duplication preserves perfection.

Proof. Suppose otherwise. Since every induced subgraph of G is either an induced subgraph of G_1 , or of G_2 , or arises from induced subgraphs of G_1, G_2 by substitution, we may assume that G is minimally imperfect. Clearly, G_2 has at least two vertices, and $G_1 \setminus v$ has at least one vertex. Take an arbitrary vertex u of G_2 , and denote by N its neighbors of G in $V(G_1 \setminus v)$. Notice that for each vertex in $V(G_2)$, its neighbors of G in $V(G_1 \setminus v)$ is precisely N. As G is minimally imperfect, \overline{G} is minimally imperfect by Corollary 5.8, so \overline{G} is connected, implying in turn that $V(G_1 \setminus v) - N \neq \emptyset$. Let $X := \{u\} \cup N$. Then X is a star cutset as u is adjacent to all the vertices in N, and in $G \setminus X$, there are no edges between $V(G_2) - \{u\}$ and $V(G_1 \setminus v) - N$. This contradicts the Star Cutset Lemma 5.10.

Next time we will see another important lemma about minimally imperfect graphs.

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