# CO 750 Packing and Covering: Lecture 24 

Ahmad Abdi

July 25, 2017

### 11.2 The sums of circuits property

Take an integer $n \geq 1$ and a set $S \subseteq\{0,1\}^{n}$. Recall that $S$ is a binary space if $\mathbf{0} \in S$, and if $a, b \in S$ then $a \triangle b \in S$. When is a binary space cube-ideal? The following gives a partial characterization of the cube-ideal binary spaces:

Theorem 11.5 (Abdi, Cornuéjols, Guričanová, Lee 2018+). Take an integer $n \geq 1$ and a binary space $S \subseteq$ $\{0,1\}^{n}$. Then $S$ is cube-ideal if, and only if,

$$
\operatorname{conv}(S)=\left\{x \in[0,1]^{n}: x(F)-x(D-F) \leq|F|-1 \quad \forall \text { cocycles } D \text { and odd subsets } F \subseteq D\right\} .
$$

Proof. $(\Leftarrow)$ Notice that each inequality $x(F)-x(D-F) \leq|F|-1$ can be rewritten as

$$
\sum_{i \in D-F} x_{i}+\sum_{j \in F}\left(1-x_{j}\right) \geq 1,
$$

which is a generalized set covering inequality. Thus $S$ is cube-ideal. $(\Rightarrow)$ Suppose coversely that $S$ is cube-ideal. We first prove that

$$
\operatorname{conv}(S) \subseteq\left\{x \in[0,1]^{n}: x(F)-x(D-F) \leq|F|-1 \quad \forall \text { cocycles } D \text { and odd subsets } F \subseteq D\right\} .
$$

Denote by $P$ the polytope on the right. To prove this inclusion, it suffices to show that for every cycle $C, \chi_{C}$ belongs to $P$. Well, for every cocycle $D$ and odd subset $F \subseteq D$, we have $C \cap D \neq F$ because $|C \cap D|$ is even, implying in turn that

$$
\chi_{C}(F)-\chi_{C}(D-F) \leq 1 .
$$

Thus, $\chi_{C} \in P$. To prove the reverse inclusion, it suffices to prove that every inequality defining $\operatorname{conv}(S)$ is valid for $P$. Since $S$ is cube-ideal, $\operatorname{conv}(S)$ is described by hypercube inequalities - which are valid for $P-$ and by generalized set covering inequalities. Take disjoint subsets $I, J \subseteq[n]$ such that $\sum_{i \in I} x_{i}+\sum_{j \in J}\left(1-x_{j}\right) \geq 1$ is a defining inequality of $\operatorname{conv}(S)$.

Claim. There is a cocycle $D$ such that $D \subseteq I \cup J$ and $|D \cap J|$ is odd.

Proof of Claim. To see this, write

$$
S=\left\{x \in\{0,1\}^{[n]}: A x \equiv \mathbf{0} \quad(\bmod 2)\right\}
$$

for some $0-1$ matrix $A$. Let $d$ be the sum of the columns in $J$ of $A$, and let $B$ be the submatrix of $A$ obtained after dropping columns $I \cup J$. Since $\sum_{i \in I} x_{i}+\sum_{j \in J}\left(1-x_{j}\right) \geq 1$ is valid for every point of $S$, the binary system

$$
B y \equiv d \quad(\bmod 2)
$$

has no $0-1$ solution. (For if $y$ is a solution, then by setting the coordinates in $I$ to 0 and the coordinates in $J$ to 1 , we can extend $y$ to a cycle $x$ for which $\sum_{i \in I} x_{i}+\sum_{j \in J}\left(1-x_{j}\right)=0$, which is not the case.) By Farkas' lemma for binary spaces, there is a $0-1$ vector $c$ such that $c^{\top} B \equiv \mathbf{0}(\bmod 2)$ and $c^{\top} d \equiv 1(\bmod 2)$. Consider the cocycle $\chi_{D}:=c^{\top} A$. Then the first equation implies that $D \subseteq I \cup J$, while the second equation implies that $|D \cap J|$ is odd, as required.

Let $F:=D \cap J$. Then $F$ is an odd subset of the cocycle $D$. Observe that the inequality

$$
\sum_{i \in I} x_{i}+\sum_{j \in J}\left(1-x_{j}\right) \geq 1
$$

is dominated by the inequality

$$
\sum_{i \in D-F} x_{i}+\sum_{j \in F}\left(1-x_{j}\right) \geq 1 \quad \text { which is equivalent to } \quad x(F)-x(D-F) \leq|F|-1
$$

because $D-F \subseteq I$ and $F \subseteq J$. As a result, every inequality defining $\operatorname{conv}(S)$ is valid for $P$, so $\operatorname{conv}(S) \supseteq P$. Hence, $\operatorname{conv}(S)=P$, thereby finishing the proof.

Consider the polyhedral cone generated by $S$ :

$$
\operatorname{cone}(S)=\left\{\sum_{x \in S} \alpha_{x} x: \alpha \in \mathbb{R}_{+}^{S}\right\} \subseteq\{0,1\}^{n}
$$

Barahona and Grötschel (1986) showed that due to the transitivity of $S$, to describe the facets of $\operatorname{conv}(S)$, it suffices to have a facet description of cone $(S)$.

Theorem 11.6 (Barahona and Grötschel 1986). Take an integer $n \geq 1$ and a binary space $S \subseteq\{0,1\}^{n}$. Then

$$
\begin{equation*}
\operatorname{conv}(S)=\left\{x \in[0,1]^{n}: x(F)-x(D-F) \leq|F|-1 \quad \forall \text { cocycles } D \text { and odd subsets } F \subseteq D\right\} \tag{1}
\end{equation*}
$$

if, and only if,

$$
\begin{equation*}
\operatorname{cone}(S)=\left\{x \in \mathbb{R}_{+}^{n}: x_{f}-x(D-\{f\}) \leq 0 \quad \forall \text { cocycles } D \text { and } f \in D\right\} \tag{2}
\end{equation*}
$$

Proof. $(\Rightarrow)$ Suppose that (1) holds. As $\mathbf{0} \in S$, the facets of $\operatorname{conv}(S)$ tight at $\mathbf{0}$ describe the conic hull of $S$. Since the cocycle inequality

$$
x(F)-x(D-F) \leq|F|-1 \quad \text { cocycle } D \text { and odd subset } F \subseteq D
$$

is tight at $\mathbf{0}$ if and only if $|F|=1$, it follows that (2) holds. $(\Leftarrow)$ Conversely, suppose that (2) holds. To prove that (1) holds, let

$$
(\diamond) \quad \sum_{i \in I} a_{i} x_{i}+\sum_{j \in[n]-I} a_{j}\left(1-x_{j}\right) \geq b \quad a \in \mathbb{R}_{+}^{n}, b \in \mathbb{R}
$$

be a facet-defining inequality for $\operatorname{conv}(S)$. It suffices to show that $(\diamond)$ is equivalent to a cocycle inequality. To this end, take a point $u \in S$ that lies on this facet. Consider the change of variables $x_{i} \mapsto 1-x_{i}$ for the indices in $\left\{i \in[n]: u_{i}=1\right\}$; this mapping sends the above inequality to the inequality

$$
\sum_{i \in I: u_{i}=0} a_{i} x_{i}+\sum_{i \in I: u_{i}=1} a_{i}\left(1-x_{i}\right)+\sum_{j \in[n]-I: u_{j}=0} a_{j}\left(1-x_{j}\right)+\sum_{j \in[n]-I: u_{j}=1} a_{j} x_{j} \geq b
$$

and the set $S$ to the set $S \triangle u:=\{x \Delta u: x \in S\}$. Then $(\star)$ is a facet-defining inequality for $S \triangle u$ and the facet contains the point $\mathbf{0}=u \triangle u \in S \triangle u$. Hence, $(\star)$ also defines a facet for cone $(S \triangle u)$. However, since $S$ is a binary space, $S \triangle u$ is just the original set $S$, so $(\star)$ defines a facet of cone $(S)$. By (2), there is a cocycle $D \subseteq[n]$ and an element $f \in D$ such that $(\star)$ is equivalent to the inequality

$$
x_{f}-x(D-\{f\}) \leq 0
$$

Take the cycle $C \subseteq[n]$ such that $u=\chi_{C}$. Reverting back the change of variables, we see that $(\diamond)$ is equivalent to

$$
x(F)-x(D-F) \leq|F|-1
$$

where $F=(C \cap D) \triangle\{f\}$. Since $|C \cap D|$ is even, it follows that $|F|$ is odd, so $(\diamond)$ is equivalent to a cocycle inequality, as required.

Take a finite set $E$ and a binary space $S \subseteq\{0,1\}^{E}$. A circuit is a non-empty cycle that does not properly contain another non-empty cycle. Notice that $\left\{\chi_{C}: C\right.$ is a circuit $\}$ are precisely the non-zero points in $S$ of minimal support.

Remark 11.7. Take a finite set $E$ and a binary space $S \subseteq\{0,1\}^{E}$, and take a non-empty subset $C \subseteq E$. Then $C$ is a cycle if, and only if, $C$ is a disjoint union of circuits.

Proof. $(\Leftarrow)$ If $C$ is a disjoint union of circuits, then it is also a symmetric difference of circuits, so $C$ is a cycle. $(\Rightarrow)$ We prove the converse by induction on $|C|$. As cycles of length 1 are already circuits, the base case $|C|=1$ holds. For the induction step, assume that $|C| \geq 2$. If $C$ does not properly contain a non-empty cycle, then it is already a circuit, and we are done. Otherwise, there is a non-empty cycle $C^{\prime}$ such that $C^{\prime} \subsetneq C$. Let $C^{\prime \prime}:=C \triangle C^{\prime}$. Notice that $C$ is the disjoint union of the non-empty cycles $C^{\prime}, C^{\prime \prime}$. By the induction hypothesis, each of $C^{\prime}, C^{\prime \prime}$ is the disjoint union of circuits, implying in turn that $C$ is a disjoint union of circuits, thereby completing the induction step.

It follows from Remark 11.7 that cone $(S)$ is the polyhedral cone generated by the circuits, that is,

$$
\operatorname{cone}(S)=\left\{\sum_{C \text { is a circuit }} y_{C} \cdot \chi_{C}: y \in \mathbb{R}_{+}^{\text {circuits }}\right\} .
$$

The binary space $S$ has the sums of circuits property if

$$
\left\{\sum_{C \text { is a circuit }} y_{C} \cdot \chi_{C}: y \in \mathbb{R}_{+}^{\text {circuits }}\right\}=\left\{x \in \mathbb{R}_{+}^{n}: x_{f}-x(D-\{f\}) \leq 0 \quad \forall \text { cocycles } D \text { and } f \in D\right\} .
$$

Equivalently, a binary space $S$ has the sums of circuits property if for each $w \in \mathbb{R}_{+}^{E}$ satisfying

$$
w(D-\{f\}) \geq w_{f} \quad \forall \text { cocycles } D \text { and } f \in D
$$

there is a vector $y \in \mathbb{R}_{+}^{\text {circuits }}$ such that $w=\sum\left(y_{C} \cdot \chi_{C}: C\right.$ is a circuit $)$. Combining Theorems 11.5 and 11.6, we get the following:

Corollary 11.8. A binary space is cube-ideal if, and only if, it has the sums of circuits property.

The deep concept of the sums of circuits property was introduced by Seymour (1979), where he showed that the binary space corresponding to the cycles of a graph has the sums of circuits property. (This result is in fact equivalent to Corollary 8.12.) That is, given a graph $G=(V, E)$, if a vector $w \in \mathbb{Z}_{+}^{E}$ satisfies

$$
w(D-\{f\}) \geq w_{f} \quad \forall \text { cuts } D \text { and } f \in D
$$

then there is an assignment $y \in \mathbb{R}_{+}^{\text {circuits }}$ such that

$$
w=\sum\left(y_{C} \cdot \chi_{C}: C \text { is a circuit }\right) .
$$

In particular, if $G=(V, E)$ is a bridgeless graph, then there is an assignment $y \in \mathbb{R}_{+}^{\text {circuits }}$ such that

$$
\mathbf{1}=\sum\left(y_{C} \cdot \chi_{C}: C \text { is a circuit }\right)
$$

Szekeres (1973), and independently Seymour (1979), conjecture that $y$ can be chosen to be half-integral; this is known as the notoriously difficult cycle double cover conjecture.

