

MA431 Lecture 6

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Outline

- 1 Recap
- 2 Flows of minimum energy
- 3 Projection onto cut space
- 4 Electrical flows and electrical potentials
- 5 Computing i and Π
- 6 Effective resistance
- 7 Kirchhoff's effective resistance theorem
- 8 A strengthening of Kirchhoff's effective resistance theorem
- 9 Uniform sampling of spanning trees

The incidence matrix of an orientation

- Let $G = (V, E)$ be a **connected** graph with n vertices and m edges.
- Denote by $\vec{G} = (V, \vec{E})$ an arbitrary orientation of G .

Recall

The **incidence matrix** of \vec{G} is the $V \times \vec{E}$ matrix B where column $(v, u) \in \vec{E}$ is equal to $e_u - e_v$.

Recall

The Laplacian matrix of G is BB^T .

Cycle and cut spaces

Cycle space

$$W^\diamond = \left\{ f \in \mathbb{R}^{\vec{E}} : Bf = \mathbf{0} \right\}$$

Cut space

$$W^\star = \left\{ B^\top \pi : \pi \in \mathbb{R}^V \right\}$$

Remark

- 1 $W^\star = (W^\diamond)^\perp$ and $W^\diamond = (W^\star)^\perp$
- 2 $\mathbb{R}^{\vec{E}} = W^\diamond \oplus W^\star$.

Cycle and cut spaces

T spanning tree

Theorem G connected

$$\dim(W^\diamond) = m - n + 1 \text{ and } \dim(W^*) = n - 1.$$

Consequence for the cycle space

- 1 $\{\chi(C_e) : e \in \vec{E} \setminus T\}$ is a basis for W^\diamond .
- 2 $W^\diamond = \text{span}\{\chi(C) : C \text{ is a directed cycle of } \vec{G}\}$.

Consequence for the cut space

- 1 $\{\chi(\delta^+(S_e)) : e \in T\}$ is a basis for W^* .
- 2 $W^* = \text{span}\{\chi(\delta^+(S)) : S \subseteq V, S \neq \emptyset, V\}$.
- 3 $W^* = \text{span}\{\chi(\delta^+(r)) : r \in V\}$.
- 4 $\{\chi(\delta^+(r)) : r \in V \setminus t\}$ is a basis for W^* , for any $t \in V$.

∇ , grad, and ∇grad

The linear operator ∇

Let $f \in \mathbb{R}^{\vec{E}}$. Think of f as a **flow** on \vec{G} where $f_{(u,v)}$ is the amount of flow going **from u to v** .

Net flow

For each $v \in V$, define

$$\nabla f_v = (\nabla f)_v := \sum_{(u,v) \in \vec{E}} f_{(u,v)} - \sum_{(v,w) \in \vec{E}} f_{(v,w)}.$$



$\nabla : \mathbb{R}^{\vec{E}} \rightarrow \mathbb{R}^V$ is a linear operator represented by the matrix B . That is, for each $f \in \mathbb{R}^{\vec{E}}$, the image ∇f is the matrix product Bf .

Definition

A **circulation** of \vec{G} is a flow $f \in \mathbb{R}^{\vec{E}}$ s.t. $\nabla f = \mathbf{0}$.

Remark

W^\diamond is precisely the set of all circulations of \vec{G} .

The linear operator grad

Definition (gradient)

Let $\text{grad} : \mathbb{R}^V \rightarrow \mathbb{R}^{\vec{E}}$ be the operator defined as follows: for each $\pi \in \mathbb{R}^V$,

$$(\text{grad } \pi)_{(u,v)} = \pi_v - \pi_u$$



grad is a linear operator represented by the matrix B^T . That is, for each $\pi \in \mathbb{R}^V$, the image $\text{grad } \pi$ is the matrix product $B^T \pi$.

Range of grad

$$\text{range}(\text{grad}) = \{B^T \pi : \pi \in \mathbb{R}^V\} = \omega^*$$

The Laplacian as a differential operator

Summary

- 1 ∇ is represented by B .
- 2 grad is represented by B^T .

Lemma

∇ and grad are adjoint operators: $\langle \nabla f, \pi \rangle = \langle f, \text{grad } \pi \rangle$.

Lemma

$\nabla \text{grad} : \mathbb{R}^V \rightarrow \mathbb{R}^V$ is a linear operator represented by the Laplacian matrix L of G .

$$L = BB^T$$

Electrical flows

Energy

Let $b \in \mathbb{R}^V$ be a demand vector such that $b = \nabla f$ for some flow $f \in \mathbb{R}^{\vec{E}}$.

Energy

The **energy** of f is $\|f\|^2 = \sum_{e \in \vec{E}} f_e^2$.

Goal

Find a flow of **minimum energy** satisfying the demands b . That is, solve the quadratic program:

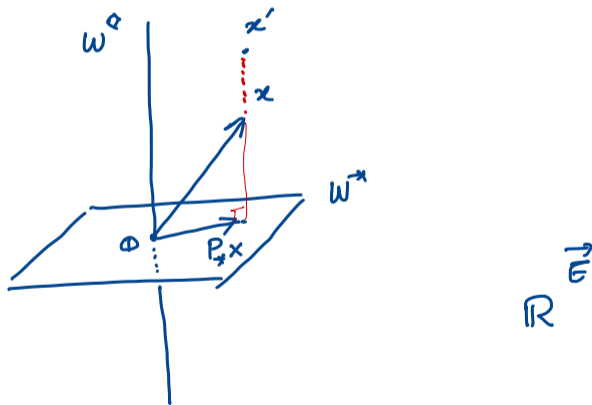
$$\min \left\{ \|f'\|^2 : \nabla f' = b, f' \in \mathbb{R}^{\vec{E}} \right\}$$

Answer: The flow turns out to be unique!
Let's find it.

Orthogonal projection onto cut space

Definition

Denote by $P_{\star} : \mathbb{R}^{\vec{E}} \rightarrow \mathbb{R}^{\vec{E}}$ the linear operator that projects orthogonally onto W^{\star} .



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Lemma

If $\nabla f' = \nabla f$, then $P_{\star} f' = P_{\star} f$.

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Lemma

If $\nabla f' = \nabla f$, then $P_{\star} f' = P_{\star} f$.

Proof.

- $\nabla(f' - f) = \nabla f' - \nabla f = \mathbf{0}$, so $f' - f \in W^{\diamond}$.
- Thus, f', f have the same orthogonal projection onto $(W^{\diamond})^{\perp} = W^{\star}$.



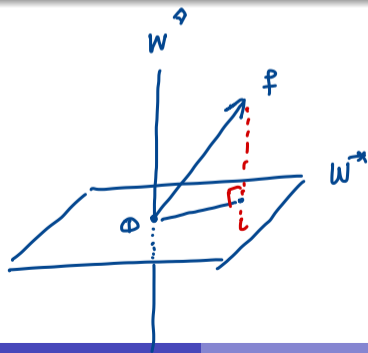
Flow of minimum energy

Let $b \in \mathbb{R}^V$ be a demand vector such that $b = \nabla f$ for some flow $f \in \mathbb{R}^{\vec{E}}$.

Theorem

Let $\iota := P_{\star} f$. Then ι is the unique optimal solution to

$$\min \left\{ \|f'\|^2 : \nabla f' = b, f' \in \mathbb{R}^{\vec{E}} \right\}$$



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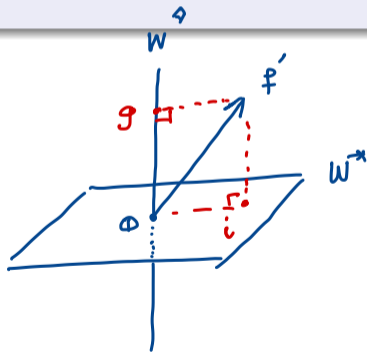
Let $\iota := P_{\star} f$. Then ι is the unique optimal solution to

$$\min \left\{ \|f'\|^2 : \nabla f' = b, f' \in \mathbb{R}^{\vec{E}} \right\}$$

Proof.

- Let f' be an arbitrary feasible solution.
- We know $P_{\star} f' = \iota$.
- Let $g \in \mathbb{R}^{\vec{E}}$ be the orthogonal projection of f' onto W^{\diamond} .
- Then

$$f' = g + i$$
$$\|f'\|^2 = \|g\|^2 + \|i\|^2$$



Flow of minimum energy

Claim 1

$\nabla \iota = b$. That is, ι is also a feasible solution to the quadratic program.

$$\nabla \iota = \nabla (f' - g') = \nabla f' - \underbrace{\nabla g'}_0 = \nabla f = b$$

Flow of minimum energy

Claim 1

$\nabla \iota = b$. That is, ι is also a feasible solution to the quadratic program.

g is a circulation so $\nabla g = \mathbf{0}$, so $\nabla \iota = \nabla f - \nabla g = b$.

Claim 2

$\|f'\|^2 \geq \|\iota\|^2$ with equality iff $f' = \iota$.

Follows from $\|f'\|^2 = \|g\|^2 + \|\iota\|^2$.

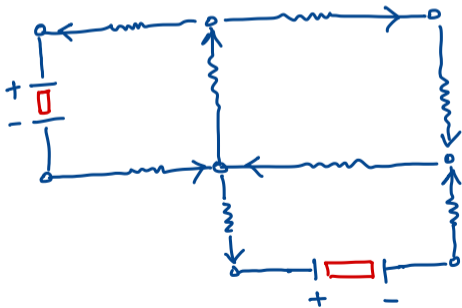


Electrical flow

Definition

ι is called the **electrical flow** satisfying the demands b .

Think of \vec{G} as an **electrical network**, the vertices as **nodes**, and the arcs as **resistors** of unit resistance, and therefore unit conductance, connecting the end nodes.



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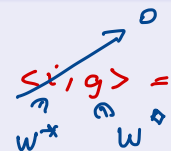
Think of \vec{G} as an **electrical network**, the vertices as **nodes**, and the arcs as **resistors** of unit resistance, and therefore unit conductance, connecting the end nodes.

Lemma

$$\langle \iota, f \rangle = \|\iota\|^2.$$

Proof.

Recall $f = i + g$ for some $g \in W^\diamond$. Then

$$\langle \iota, f \rangle = \langle \iota, i + g \rangle = \langle \iota, i \rangle + \underbrace{\langle \iota, g \rangle}_{\substack{= \\ \underbrace{\langle \iota, g \rangle}_{g \in W^\diamond}}} = \langle \iota, i \rangle$$


□

Electrical potentials

We know $\iota \in W^* = \text{range}(\text{grad})$. Thus, $\iota = \text{grad } \pi$ for some $\pi \in \mathbb{R}^V$.



$$i_e = \pi_u - \pi_v \quad e = (u, v)$$

Electrical potentials

We know $\iota \in W^* = \text{range}(\text{grad})$. Thus, $\iota = \text{grad } \pi$ for some $\pi \in \mathbb{R}^V$.

Definition

$\pi_v, v \in V$ are called the **electrical potentials**.

Kirchhoff's second law (potential difference)

$\iota_{(u,v)} = \pi_v - \pi_u$ for all $(u, v) \in \vec{E}$.

Electrical potentials

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Definition

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Kirchhoff's second law (potential difference)

$\iota_{(u,v)} = \pi_v - \pi_u$ for all $(u, v) \in \vec{E}$.

Remark

The electrical potentials are unique up to shifting by a constant.

$$\text{grad} : \mathbb{R}^V \rightarrow \mathbb{R}^{\vec{E}}$$

Matrix of projection onto W^*

Recall

$P_\star : \mathbb{R}^{\vec{E}} \rightarrow \mathbb{R}^{\vec{E}}$ is the linear operator that projects orthogonally onto W^* .

Let Π be the $\vec{E} \times \vec{E}$ matrix where column e is precisely $P_\star \chi^e$.

Remark

Π represents the linear operator P_\star in the basis $\{\chi^e : e \in \vec{E}\}$.

Matrix of projection onto W^*

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Π represents the linear operator P_\star in the basis $\{\chi^e : e \in \vec{E}\}$.

Remark

Π is a ^{orthogonal} projection matrix, so $\Pi^2 = \Pi$ and $\Pi^\top = \Pi$.

idempotent

self-adjoint

$$\langle x, \Pi y \rangle = \langle \Pi x, y \rangle$$

Computing ι and Π

Question

Given $b = \nabla f$, how do we compute ι ?

- 1 We know $\iota_e = \pi_v - \pi_u$ for each $e = (u, v) \in \vec{E}$. That is, $\iota = \text{grad } \pi$.
- 2 So how do we compute π ?
- 3 We know

$$b = \nabla \iota = \nabla \text{grad } \pi = L \pi$$

- 4 We can get one solution π by solving

$$\pi = L^+ b$$

L^+ : pseudoinverse of L

- 5 Putting it altogether we get

$$\begin{aligned} \Pi f &= \iota = \text{grad } \pi = \text{grad } L^+ b = \text{grad } L^+ \nabla f \\ &= B^T L^+ B f \end{aligned}$$

- 6 Then

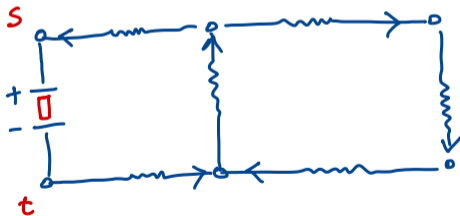
$$\Pi = B^T L^+ B$$

Effective resistance between two nodes

Take distinct vertices $s, t \in V$. Send an electrical unit flow ι' from s to t , with electrical potentials $\pi'_v, v \in V$.

Definition

The potential difference $\pi'_t - \pi'_s$ is the **effective resistance between s and t** .

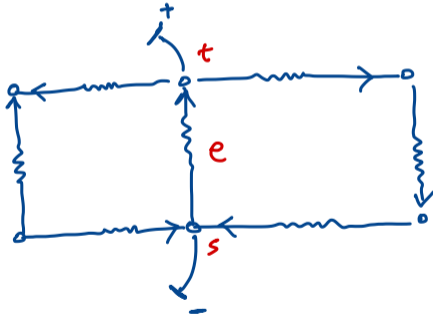


Effective resistance of an edge

Take an edge $e = (s, t) \in \vec{E}$.

Definition

The **effective resistance of edge e** is the effective resistance between s and t .



Effective resistance of an edge



Take an edge $e = (s, t) \in \vec{E}$.

Definition

The **effective resistance of edge e** is the effective resistance between s and t .

Theorem (Characterization of Effective Resistance)

Let ι' be the electrical unit flow from s to t . Let $\pi'_v, v \in V$ be the electrical potentials. Then

- 1 $\pi'_t - \pi'_s = \iota'_e$ by Kirchoff's 2nd law
- 2 the effective resistance of e is ι'_e

Effective resistance of an edge

Take an edge $e = (s, t) \in \vec{E}$.

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The **effective resistance of edge e** is the effective resistance between s and t .

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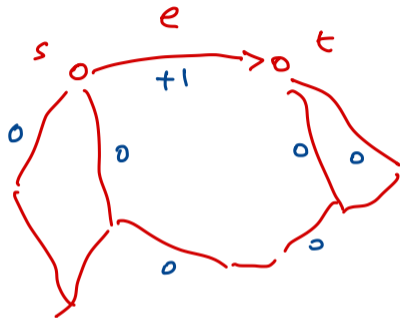
- 1 $\pi'_t - \pi'_s = \iota'_e$
- 2 the effective resistance of e is ι'_e
- 3 the effective resistance of e is $\Pi_{e,e}$
- 4 $\iota'_e = \|\iota'\|^2$.

Proof of Theorem

x^e

$$\textcircled{3} \quad \Pi_{e,e} = (P_x x^e)_e$$

Note that x^e is a flow of unit value from s to t !



$$= (i')_e$$
$$= i'_e$$

$$\textcircled{4} \quad \|i'\|^2 = \langle i', i' \rangle = \langle i', x^e \rangle = i'_e .$$

another flow satisfying the same demands,

Effective resistance and connectivity

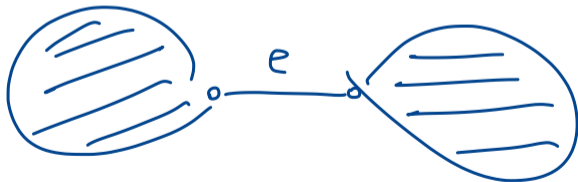
Lemma

- 1 The effective resistance of $e \in \vec{E}$ is at most 1.
- 2 The effective resistance of $e \in \vec{E}$ is 1 if and only if $G \setminus e$ is not connected.

Proof.

Exercise. □

②



Intuition: The higher the effective resistance, the "more important" the edge.

Effective resistance and connectivity

Lemma

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- 2 The effective resistance of $e \in \vec{E}$ is 1 if and only if $G \setminus e$ is not connected.

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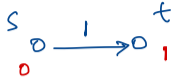
Exercise. □

Rayleigh monotonicity principle

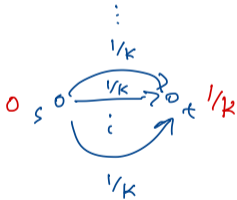
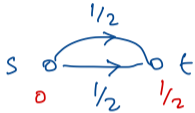
Let $G' = (V, E')$ with $E' \supseteq E$. Then for any two vertices $u, v \in V$, the effective resistance between u and v in G' is smaller than or equal to the effective resistance between u and v in G .

Proof.

Exercise. □



π'
 i'



k edges

Kirchhoff's effective resistance theorem

Take an edge $e = (s, t) \in \vec{E}$.

Recall

$T(G)$ is the number of spanning trees of G .

Theorem

The effective resistance of e is $\frac{T(G/e)}{T(G)}$.

$T(G/e) = \#$ spanning trees of G that contain e

$\therefore \frac{T(G/e)}{T(G)} =$ the proportion of spanning trees that contain e

Kirchhoff's effective resistance theorem

Take an edge $e = (s, t) \in \vec{E}$.

Recall

$T(G)$ is the number of spanning trees of G .

Theorem

The effective resistance of e is $\frac{T(G/e)}{T(G)}$.

Relabel $V = \{1, \dots, n\}$ and $e = (n, n-1)$.



Recall

Let ι' be the electrical unit flow from n to $n-1$. Let $\pi'_v, v \in V$ be the electrical potentials. Then the effective resistance of e is $\pi'_{n-1} - \pi'_n$.

Proof of Kirchhoff's effective resistance theorem

- Let us apply Cramer's rule to solve

$$L\pi' = \underbrace{e_{n-1} - e_n}_{\text{demand vector}}$$

Proof of Kirchhoff's effective resistance theorem

- Let us apply Cramer's rule to solve

$$L\pi' = e_{n-1} - e_n$$

- The rows of L add up to $\mathbf{0}$, so we may drop the last row of $L\pi' = e_{n-1} - e_n$.

. Force $\pi'_n = 0$.

Proof of Kirchhoff's effective resistance theorem

- Let us apply Cramer's rule to solve

$$L\pi' = e_{n-1} - e_n$$

- The rows of L add up to $\mathbf{0}$, so we may drop the last row of $L\pi' = e_{n-1} - e_n$.
- To ensure that π' is unique, we enforce in addition

$$\pi'_n = 0$$

The diagram illustrates the reduction of the matrix equation $L\pi' = e_{n-1} - e_n$. It shows two equivalent systems:

- On the left, a matrix $L[n]$ (with a red border) is multiplied by a vector π' (with a red border). The vector π' has components $0, 0, \dots, 0, 1, -1$. The matrix $L[n]$ has a vertical line on the right side, and the bottom row is marked with \dots .
- A tilde symbol \sim indicates the second system, where the matrix $L[n]$ (with a blue border) is multiplied by a vector π' (with a red border). The vector π' has components $0, \dots, 0, 0, 1$. The matrix $L[n]$ has a vertical line on the right side, and the bottom row is marked with $0, \dots, 0, 0, 1$.
- On the right, the vector π' (with a red border) is shown as $0, \dots, 0, 0$.

Proof of Kirchhoff's effective resistance theorem

- Let us apply Cramer's rule to solve

$$L\pi' = e_{n-1} - e_n$$

- The rows of L add up to $\mathbf{0}$, so we may drop the last row of $L\pi' = e_{n-1} - e_n$.
- To ensure that π' is unique, we enforce in addition

$$\pi'_n = 0$$

- As a result, π' is obtained by solving

$$L[n] y = e_{n-1}$$

- By Cramer's rule,

$$\pi'_{n-1} - \pi'_n = \pi'_{n-1} = y_{n-1} =$$

$$\frac{\det(L[n][n-1])}{\det(L[n])}$$

See the Matrix-Tree
Thm & its proof

$$= \frac{T(G/e)}{T(G)}$$

The diagonal entries of Π

Kirchhoff's effective resistance theorem

The effective resistance of e is $\frac{T(G/e)}{T(G)}$.

Recall

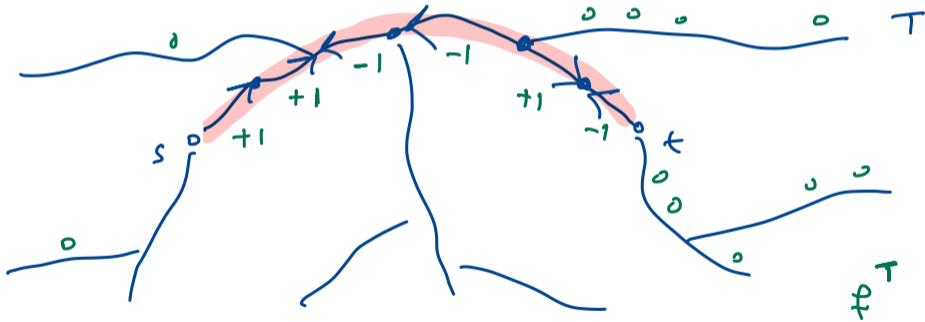
The effective resistance of e is $\Pi_{e,e}$.

Corollary

$$\Pi_{e,e} = \frac{T(G/e)}{T(G)}.$$

A strengthening of Kirchhoff's effective resistance theorem

- Take two vertices $s, t \in V$.
- For each spanning tree $T \subseteq \vec{E}$, let $f^T \in \mathbb{R}^{\vec{E}}$ be the flow that sends 1 unit of flow from s to t along the unique st -path in T .



A strengthening of Kirchhoff's effective resistance theorem

- Take two vertices $s, t \in V$.
- For each spanning tree $T \subseteq \vec{E}$, let $f^T \in \mathbb{R}^{\vec{E}}$ be the flow that sends 1 unit of flow from s to t along the unique st -path in T .
- Let

$$g := \frac{1}{T(G)} \sum_{T \in \mathcal{T}} f^T$$

where \mathcal{T} is the set of all spanning trees.

Theorem

g is the electrical unit flow from s to t .

*This implies Kirchhoff's effective resistance theorem
(set $e = (s, t)$ & look at g_e).*

Proof

Claim 1

$$\nabla g = e_t - e_s.$$

Proof

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$$\nabla g = e_t - e_s.$$

$$\nabla g = \frac{1}{T(G)} \sum_{T \in \mathcal{T}} \nabla f^T = \frac{1}{T(G)} \cdot T(G) \cdot (e_t - e_s) = e_t - e_s$$

Claim 2

$$g \in W^* = (W^\diamond)^\top.$$

- Let $C \subseteq \vec{E}$ be a directed cycle in \vec{G} .
- It suffices to show $g^\top \chi(C) = 0$
- This is invariant under the orientation of G , so WMA that C is a directed cycle in \vec{G} .

Proof

Thus,

$$g^T \chi(C) = \sum_{a \in C} g_a$$

Proof

Thus,

$$\begin{aligned}g^T \chi(C) &= \sum_{a \in C} g_a \\ &= \frac{1}{T(G)} \sum_{a \in C} \sum_{T \in \mathcal{T}} f_a^T\end{aligned}$$

Proof

Thus,

$$\begin{aligned}g^T \chi(C) &= \sum_{a \in C} g_a \\&= \frac{1}{T(G)} \sum_{a \in C} \sum_{T \in \mathcal{T}} f_a^T \\&= \frac{1}{T(G)} \sum \left(f_a^T : a \in C, T \in \mathcal{T}, a \in T \right)\end{aligned}$$

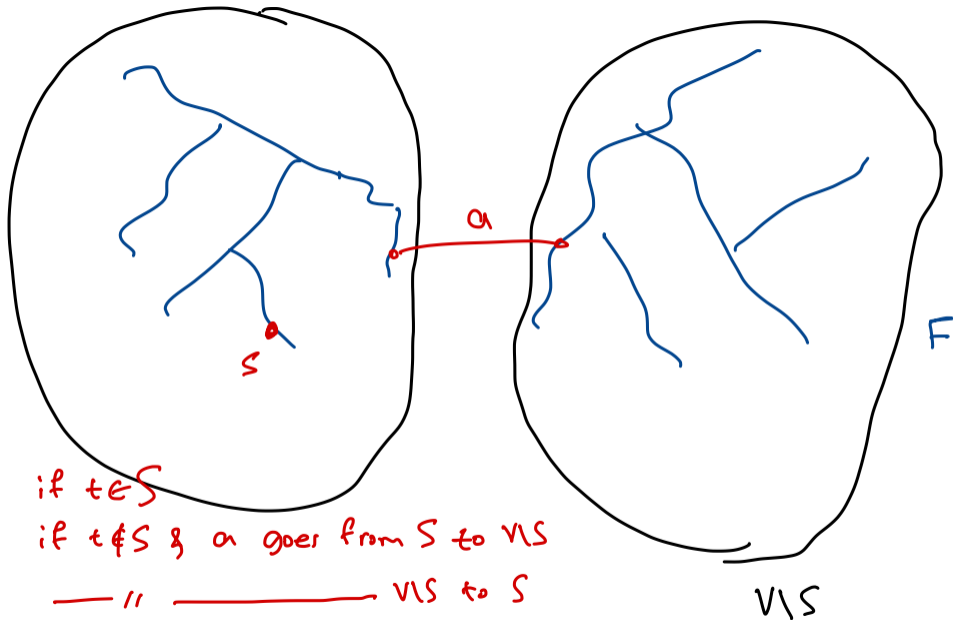
Proof

Thus,

$$\begin{aligned}g^T \chi(C) &= \sum_{a \in C} g_a \\&= \frac{1}{T(G)} \sum_{a \in C} \sum_{T \in \mathcal{T}} f_a^T \\&= \frac{1}{T(G)} \sum \left(f_a^T : a \in C, T \in \mathcal{T}, a \in T \right) \\&= \frac{1}{T(G)} \sum_F \sum \left(f_a^T : a \in C, T \in \mathcal{T}, T \setminus a = F \right)\end{aligned}$$

where $F \subseteq \vec{E}$ is any spanning forest with two connected components.

Proof



$$t_a^{\pi T} = \begin{cases} 0 & \text{if } t \in S \\ +1 & \text{if } t \notin S \text{ \& } a \text{ goes from } S \text{ to } VIS \\ -1 & \text{if } a \text{ goes from } VIS \text{ to } S \end{cases}$$

Proof

Fix F . Let $S, V \setminus S$ be the two connected components where $s \in S$.

- Case 1: $t \in S$. Then, if $T \in \mathcal{T}$ and a satisfies $T \setminus a = F$, then $f_a^T = 0$, so

$$\sum \left(f_a^T : a \in \mathcal{C}, T \in \mathcal{T}, T \setminus a = F \right) = 0$$

Proof

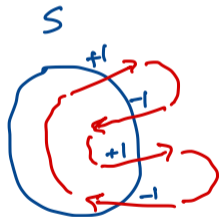
Fix F . Let $S, V \setminus S$ be the two connected components where $s \in S$.

- Case 2: $t \in V \setminus S$. Then

$$\begin{aligned} & \sum \left(f_a^T : a \in C, T \in \mathcal{T}, T \setminus a = F \right) \\ &= \sum (1 : a \in C, \text{tail}(a) \in S, \text{head}(a) \in V \setminus S) \\ & \quad + \sum (-1 : a \in C, \text{tail}(a) \in V \setminus S, \text{head}(a) \in S) \end{aligned}$$

$$= |C \cap \delta^+(S)| - |C \cap \delta^-(S)|$$

$$= 0 \quad \text{b/c } C \text{ is a directed cycle!}$$



Proof

Thus,

$$\begin{aligned}g^T \chi(C) &= \sum_{a \in C} g_a \\&= \frac{1}{T(G)} \sum_{a \in C} \sum_{T \in \mathcal{T}} f_a^T \\&= \frac{1}{T(G)} \sum \left(f_a^T : a \in C, T \in \mathcal{T}, a \in T \right) \\&= \frac{1}{T(G)} \sum_F \underbrace{\sum \left(f_a^T : a \in C, T \in \mathcal{T}, T \setminus a = F \right)}_0 \\&= 0\end{aligned}$$

Proof

Theorem

g is the electrical unit flow from s to t .

Proof.

Claim 1

$$\nabla g = e_t - e_s.$$

Claim 2

$$g \in W^* = (W^\diamond)^\top.$$

Thus, g must be **the** electrical unit flow from s to t .

The columns of Π

Take $e = (s, t) \in \vec{E}$.

Recall

$P_{\star}\chi^e$ is the electrical unit flow from s to t .

Recall

$P_{\star}\chi^e$ is column e of Π .

The columns of Π

Take $e = (s, t) \in \vec{E}$.

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Recall

$P_{\star}\chi^e$ is column e of Π .

We just showed

$\frac{1}{T(G)} \sum_{T \in \mathcal{T}} f^T$ is the electrical unit flow from s to t .

Corollary

Column e of Π is equal to $\frac{1}{T(G)} \sum_{T \in \mathcal{T}} f^T$.

Uniform sampling of spanning trees

Recall

$$\Pi_{e,e} = \frac{T(G/e)}{T(G)}.$$

Remark

$\frac{T(G/e)}{T(G)}$ is the fraction of spanning trees of G that use the edge e .

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Kirchhoff's effective resistance theorem (rephrased)

If T is a uniformly random spanning tree, then $\Pr[e \in T] = \Pi_{e,e}$.

The Transfer-Current Theorem

The Transfer-Current Theorem

If T is a uniformly random spanning tree, then for any $F \subseteq E$,

$$\Pr[F \subseteq T] = \det(\Pi_F)$$

where Π_F denotes the principal submatrix of Π indexed by F .