On the mixing set with a knapsack constraint

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Abstract

We study a substructure appearing in mixed-integer programming reformulations of chance-constrained programs with stochastic right-hand-sides over a finite discrete distribution, which we call the mixing set with a knapsack constraint. Recently, Luedtke et al. (2010) and Küçükyavuz (2012) studied valid inequalities for such sets. However, most of their results were focused on the equal probabilities case (when the knapsack constraint reduces to a cardinality constraint). In this paper, we focus on the general probabilities case (general knapsack constraint).

We characterize the valid inequalities that do not come from the knapsack polytope and use this characterization to generalize the results previously derived for the equal probabilities case. Our results allow for a deep understanding of the relationship that the set under consideration has with the knapsack polytope. Moreover, they allow us to establish benchmarks that can be used to identify when a relaxation will be useful for the considered types of reformulations of chance-constrained programs.

1 Introduction

Many optimization problems in real world applications allow to some extent a number of violated constraints, which results in a decrease in the quality of service, as well as a decrease in the cost of production. These optimization problems have been a main motive to study probabilistic (in particular, chance-constrained) programming. A difficulty when dealing with these optimization problems is that the feasible region is not necessarily convex. In this paper, we consider mixed-integer programming (MIP) reformulations of chance-constrained programs with joint probabilistic constraints in which the right-hand-side vector is random with a finite discrete distribution[1]. This model was first proposed in Sen [16], studied in Ruszczyński [15], and extended by Luedtke et. al [14] and Küçükyavuz [12]. This reformulation gives rise to a mixing-type set [10] subject to an additional knapsack constraint, which is the focus of this paper.[2]

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[1]Sometimes the finite discrete distribution given is an approximation of an unknown continuous distribution, obtained via a Sample Average Approximation (SAA) technique.

[2]Prior to our work, the main focus seems to have been on the equal probabilities case, justified by employing the SAA technique. Our more general framework of the general probabilities case allows for more sophisticated techniques such as Importance Sampling. See the recent work of Barrera et. al [6] for further details.

Formally, consider the following chance-constrained programming problem

$$
\begin{aligned}
\min \quad & c^\top x \\
\text{s.t.} \quad & \mathbb{P}(Bx \geq \xi) \geq 1 - \epsilon \\
& x \in X,
\end{aligned}
\tag{PLP}
$$

where X is a polyhedron, B is a matrix with d rows, ξ is a random variable in $\mathbb{R}^d$ with finite discrete distribution, $\epsilon \in (0,1)$, and c is an arbitrary cost vector. Let ξ take values $\xi^1, \ldots, \xi^n$ with probabilities $\pi_1, \ldots, \pi_n$, respectively. We may assume without loss of generality that $\xi^j \geq 0$, for all $j \in [n] := \{1, \ldots, n\}$ (by a mere simple linear transformation). Thus, since $\epsilon < 1$, $Bx \geq 0$ for every feasible solution x to (PLP). By definition, for each $j \in [n]$, $\pi_j > 0$ and $\sum_{j=1}^n \pi_j = 1$. We can reformulate the chance constraint in (PLP) using linear inequalities and auxiliary binary variables as follows: let $z \in \{0,1\}^n$ where $z_j = 0$ guarantees that $Bx \geq \xi^j$. Then (PLP) is equivalent to

$$
\begin{aligned}
\min \quad & c^\top x \\
\text{s.t.} \quad & y = Bx \\
& y + \xi^j z_j \geq \xi^j \quad \forall j \in [n] \\
& \sum_{j=1}^n \pi_j z_j \leq \epsilon \\
& z \in \{0,1\}^n \\
& x \in X.
\end{aligned}
\tag{CMIP}
$$

We may assume for all $j \in [n]$ that $\pi_j \leq \epsilon$. Indeed, if $\pi_j > \epsilon$ for some $j \in [n]$, then $z_j = 0$ for all feasible solutions (x, y, z) to the above system, so we may as well drop the index j. Now let

$$
\mathcal{D} := \left\{ (y, z) \in \mathbb{R}_+^d \times \{0,1\}^n : \sum_{j=1}^n \pi_j z_j \leq \epsilon; \ y + \xi^j z_j \geq \xi^j, \ \forall j \in [n] \right\}.
$$

Then (PLP) can be rewritten as

$$
\begin{aligned}
\min \quad & c^\top x \\
\text{s.t.} \quad & Bx \in \mathrm{proj}_y \mathcal{D} \\
& x \in X.
\end{aligned}
$$

This motivates us to study the set $\mathcal{D}$. For each $\kappa \in [d]$, let

$$
\mathcal{D}_\kappa := \left\{ (y_\kappa, z) \in \mathbb{R}_+ \times \{0,1\}^n : \begin{array}{l} \sum_{j=1}^n \pi_j z_j \leq \epsilon \\ y_\kappa + \xi_\kappa^j z_j \geq \xi_\kappa^j, \ \forall j \in [n] \end{array} \right\}.
\tag{1}
$$

Then observe that

$$
\mathcal{D} = \bigcap_{\kappa \in [d]} \left\{ (y, z) \in \mathbb{R}_+^d \times \{0,1\}^n : (y_\kappa, z) \in \mathcal{D}_\kappa \right\}.
$$

Therefore, instead of $\mathcal{D}$, we study the lower dimensional sets $\mathcal{D}_\kappa$.

Fix $\kappa \in [d]$ and for notational convenience, let $h_j := \xi_\kappa^j$ for each $j \in [n]$. Let $\sum_{j \in [n]} a_j z_j \leq p$ be a valid inequality for $\mathcal{D}_\kappa$ where $a \in \mathbb{R}_+^n$, $p \in \mathbb{R}_+$, $a_j \leq p$ for all $j \in [n]$, and $\sum_{j \in [n]} a_j > p$. Observe that this inequality may be the knapsack constraint $\sum_{j=1}^n \pi_j z_j \leq \epsilon$. In this paper, we focus on the set

$$
Q := \left\{ (y, z) \in \mathbb{R}_+ \times \{0,1\}^n : \sum_{j=1}^n a_j z_j \leq p; \ y + h_i z_i \geq h_i, \ \forall i \in [n] \right\}.
$$

Specifically, we try to understand the convex hull of Q. We assume without loss of generality that $h_1 \geq h_2 \geq \ldots \geq h_n \geq 0$.

Note that the assumption that $a_j \leq p$ for all $j \in [n]$ implies that Q is a full-dimensional set. (The points $(h_1 + 1, 0), (h_1, e_1), \ldots, (h_1, e_n)$ are in Q, where e_j is the jth n-dimensional unit vector.) Also, the assumption that $\sum_{j \in [n]} a_j > p$ implies that $y \geq h_n$, for all $y \in Q$. Observe that the set Q contains as a substructure the intersection of a mixing-type set, introduced by Günlük and Pochet [10], and a knapsack constraint (hence the title of our paper). Various structural properties of $conv(Q)$ were studied in [14] and [12] when the knapsack constraint $\sum_{j \in [n]} a_j z_j \leq p$ is a cardinality constraint. In [14], a characterization of all valid inequalities for $conv(Q)$ was given[3], and in both [14] and [12], explicit classes of facet-defining inequalities were introduced.

Our contributions

The main contribution of this paper is to generalize the results in [14] and [12]. In particular, we show that the extended formulation and polynomial time separation that these two papers obtain for the cardinality constrained case follow from the results that we establish for LP relaxations of the knapsack set. We also introduce a class of explicit facet-defining inequalities for $conv(Q)$, and give a necessary condition on the constant right-hand-side term for any inequality to be facet-defining. Additionally, we provide a compact extended formulation for (CMIP) similar to the one given in [12]. Finally, we present computational experiments to illustrate how our results can be used to set benchmarks to identify when a relaxation is useful and worth a deeper investigation.

Outline

In this paper, we do not make any assumptions on the knapsack constraint. In §2 we characterize the set of all valid inequalities for $conv(Q)$. In §3 we study how different relaxations of the knapsack polytope can be used to obtain relaxations of $conv(Q)$. We discuss in §4 some properties of facet-defining inequalities of $conv(Q)$ and introduce an explicit class of facet-defining inequalities. In §5 we provide a compact integer extended formulation for the whole set (CMIP) that we are interested in. Computational experiments are provided in §6 and a conclusion in §7.

2 A characterization of valid inequalities for $conv(Q)$

To start, let $\nu := \max \{k \in \mathbb{Z} : \sum_{j=1}^{k} a_j \leq p\}$. Note $0 < \nu < n$. Define for each $k \in \{0, 1, \ldots, \nu\}$ the knapsack set

$$P_k := \left\{ z \in \{0,1\}^n : \sum_{j=1}^{n} a_j z_j \leq p; z_1 = \cdots = z_k = 1 \right\},$$

---

[3]As shown in Luedtke [13], some valid inequalities for Q are also of value in the extremely general case where each mixing constraint is replaced by a set of arbitrary linear constraints.

and define $\phi : \{0, 1, \ldots, \nu\} \times \mathbb{R}^n \to \mathbb{R}$, as follows: for $0 \le k \le \nu$ and $\alpha \in \mathbb{R}^n$, let

$$\phi(k, \alpha) := \min \left\{ \sum_{j=k+1}^{n} \alpha_j z_j : z \in P_k \right\}.$$

Note $\phi(k, \alpha) \le 0$.

The following theorem characterizes the set of all valid inequalities for $conv(Q)$.

Theorem 1. *Take $(\gamma, \alpha, \beta) \in \mathbb{R} \times \mathbb{R}^n \times \mathbb{R}$. Then*

$$\gamma y + \sum_{j \in [n]} \alpha_j z_j \ge \beta \tag{2}$$

is a valid inequality for $conv(Q)$ if, and only if, $\gamma \ge 0$ and

$$\gamma h_{k+1} + \sum_{j=1}^{k} \alpha_j + \phi(k, \alpha) \ge \beta, \quad \forall \, 0 \le k \le \nu. \tag{3}$$

Proof. Suppose that (2) is a valid inequality of $conv(Q)$ for some $(\gamma, \alpha, \beta) \in \mathbb{R} \times \mathbb{R}^n \times \mathbb{R}$. Since $(1, \mathbf{0})$ is in the recession cone of $conv(Q)$, it follows that $\gamma \ge 0$. Take $0 \le k \le \nu$. Choose $z^* \in P_k$ so that $\sum_{j=k+1}^{n} \alpha_j z_j^* = \phi(k, \alpha)$. Let $y^* = h_{k+1}$. Observe that $(y^*, z^*) \in Q$ as $z_1^* = \ldots = z_k^* = 1$. Hence, since (2) is valid for Q, it follows that

$$\beta \le \gamma y^* + \sum_{j \in [n]} \alpha_j z_j^*$$

$$= \gamma h_{k+1} + \sum_{j=1}^{k} \alpha_j + \sum_{j=k+1}^{n} \alpha_j z_j^*$$

$$= \gamma h_{k+1} + \sum_{j=1}^{k} \alpha_j + \phi(k, \alpha).$$

Since this holds for all $0 \le k \le \nu$, it follows that (3) holds.

Conversely, suppose that $\gamma \ge 0$ and (3) holds. Let $(y^*, z^*) \in Q$. Let $h_0 := +\infty$. Since $\sum_{j=1}^{n} a_j z_j^* \le p$ there exists $0 \le k \le \nu$ for which $z_1^* = \ldots = z_k^* = 1$ but $z_{k+1}^* = 0$. Observe that $z^* \in P_k$. Since $z_{k+1}^* = 0$ we get that $y^* \ge h_{k+1}$, and so as $\gamma \ge 0$ it follows that

$$\gamma y^* + \sum_{j \in [n]} \alpha_j z_j^* \ge \gamma h_{k+1} + \sum_{j=1}^{k} \alpha_j + \sum_{j=k+1}^{n} \alpha_j z_j^*$$

$$\ge \gamma h_{k+1} + \sum_{j=1}^{k} \alpha_j + \phi(k, \alpha)$$

$$\ge \beta$$

as (3) holds. Therefore, (2) is valid for (y^*, z^*), and since this is true for all $(y^*, z^*) \in Q$, (2) is a valid inequality for $conv(Q)$. $\square$ $\square$

The following proposition states that $conv(Q)$ has all the facets of the knapsack polytope as a subset of its facets.

Proposition 2. *The inequality*

$$\sum_{j \in [n]} \alpha_j z_j \geq \beta \tag{4}$$

defines a facet of $conv(Q)$ if and only if it defines a facet of $conv(P_0)$

Proof. The proof of this proposition is straightforward from Theorem 1 and the following facts: (i) any valid inequality for $conv(P_0)$ is also valid for $conv(Q)$ and conversely (ii) any inequality of the form (4) valid for $conv(Q)$ is also valid for $conv(P_0)$. $\qquad \square$ $\qquad \square$

We end this section by collecting the coefficient vectors of all valid inequalities of $conv(Q)$ as follows:

$$C := \left\{ (\gamma, \alpha, \beta) \in \mathbb{R}_+ \times \mathbb{R}^n \times \mathbb{R} : \gamma h_{k+1} + \sum_{j=1}^{k} \alpha_j + \phi(k, \alpha) \geq \beta, \forall 0 \leq k \leq \nu \right\}.$$

Following the terminology of [10], we refer to C as the *coefficient polyhedron of Q*. It is worth mentioning that the coefficient polyhedron C is reminiscent of the polar of a polyhedron.

3   A relaxation scheme for $conv(Q)$

The results in the previous section suggest that in order for us to describe $conv(Q)$, we must be able to obtain the convex hull of the knapsack polytope. Since this is a far-stretched task in itself, we consider in this section what happens if we have a polyhedral relaxation of the knapsack polytope.

For each $0 \leq k \leq \nu$, let R_k be a (strengthened) LP relaxation of P_k in the following sense: for some matrix A^k and column vector b^k we have $R_k = \{z \in [0,1]^n : A^k z \leq b^k\}$, $R_k \cap \{0,1\}^n = P_k$, and

$$R_k \subseteq \left\{ z \in [0,1]^n : \sum_{j=1}^{n} a_j z_j \leq p; z_1 = \cdots = z_k = 1 \right\}. \tag{5}$$

For each $0 \leq k \leq \nu$, define $\varphi : \{0, 1, \ldots, \nu\} \times \mathbb{R}^n \to \mathbb{R}$ as follows:

$$\varphi(k, \alpha) := \min \left\{ \sum_{j=k+1}^{n} \alpha_j z_j : z \in R_k \right\}.$$

Observe that φ is a lower bound on ϕ as $P_k \subseteq R_k$. Let

$$C_\varphi := \left\{ (\gamma, \alpha, \beta) \in \mathbb{R}_+ \times \mathbb{R}^n \times \mathbb{R} : \gamma h_{k+1} + \sum_{j=1}^{k} \alpha_j + \varphi(k, \alpha) \geq \beta, \forall\, 0 \leq k \leq \nu \right\}$$

and

$$Q_\varphi := \left\{ (y, z) \in \mathbb{R}_+ \times [0,1]^n : \gamma y + \sum_{j=1}^{n} \alpha_j z_j \geq \beta, \forall\, (\gamma, \alpha, \beta) \in C_\varphi \right\}.$$

5

Notice that $C_\varphi \subseteq C$ and consequently, $conv(Q) \subseteq Q_\varphi$. In words, what we are doing is picking polyhedral relaxations $R_0, \ldots, R_\nu$ of the knapsack sets $P_0, \ldots, P_\nu$ and defining a relaxation Q_φ of $conv(Q)$ based on them, which is argued in the following proposition.

Proposition 3. Q_φ *is an LP relaxation of* Q, *i.e.* $Q_\varphi \cap \left(\mathbb{R}_+ \times \{0,1\}^n\right) = Q$ *and*

$$Q_\varphi \subseteq \left\{ (y,z) \in \mathbb{R}_+ \times [0,1]^n : \sum_{j=1}^n a_j z_j \leq p;\ y + h_i z_i \geq h_i,\ \forall i \in [n] \right\}.$$

Proof. We first prove that $\sum_{j=1}^n a_j z_j \leq p$ is a valid inequality for Q_φ. It suffices to show $(0, -a, -p) \in C_\varphi$. Take $0 \leq k \leq \nu$. Then, as (5) holds, it follows that

$$0 \cdot h_{k+1} + \sum_{j=1}^k (-a_j) + \varphi(k, -a) \geq -\sum_{j=1}^k a_j + \left(\sum_{j=1}^k a_j - p \right) = -p.$$

As this is true for all $0 \leq k \leq \nu$, it follows that $(0, -a, -p) \in C_\varphi$ and so $\sum_{j=1}^n a_j z_j \leq p$ is a valid inequality for Q_φ.

Given $i \in [n]$ we next prove $y + h_i z_i \geq h_i$ is valid for Q_φ. It suffices to show $(\gamma, \alpha, \beta) := (1, h_i e_i, h_i) \in C_\varphi$, where e_i is the ith unit vector in $\mathbb{R}^n$. Choose $0 \leq k \leq \nu$. Observe that $\varphi(k, h_i e_i) \geq 0$, and in fact, $\varphi(k, h_i e_i) = 0$ since $\sum_{j=1}^k e_j \in P_k \subseteq R_k$. If $k < i$ then

$$\gamma h_{k+1} + \sum_{j=1}^k \alpha_j + \varphi(k, \alpha) = h_{k+1} \geq h_i = \beta.$$

Otherwise, $k \geq i$ and so

$$\gamma h_{k+1} + \sum_{j=1}^k \alpha_j + \varphi(k, \alpha) = h_{k+1} + h_i \geq h_i = \beta.$$

Hence, $(1, h_i e_i, h_i) \in C_\varphi$ implying that $y + h_i z_i \geq h_i$.

As a result,

$$Q_\varphi \subseteq \left\{ (y,z) \in \mathbb{R}_+ \times [0,1]^n : \sum_{j=1}^n a_j z_j \leq p;\ y + h_i z_i \geq h_i,\ \forall i \in [n] \right\}.$$

It remains to show $Q_\varphi \cap \left(\mathbb{R}_+ \times \{0,1\}^n \right) = Q$. The inclusion relation above implies that $Q_\varphi \cap \left(\mathbb{R}_+ \times \{0,1\}^n \right) \subseteq Q$, and since $Q \subseteq Q_\varphi$, it follows that $Q_\varphi \cap \left(\mathbb{R}_+ \times \{0,1\}^n \right) \supseteq Q$. Thus, $Q_\varphi \cap \left(\mathbb{R}_+ \times \{0,1\}^n \right) = Q$, finishing the proof. $\qquad\qquad \square \qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad \square$

3.1 Compact extended formulations for Q_φ and C_φ

Having shown that we can define a relaxation of $conv(Q)$ based on LP relaxations of the knapsack polytope, we now start with a theorem that proposes an alternate formulation of Q_φ. It provides insight into the structure of Q_φ in terms of disjunctive programming and allows the derivation of the extended formulations discussed in this subsection.

Theorem 4. *For each $0 \leq k \leq \nu$, let $S_k := \{(h_{k+1}, z) : z \in R_k\}$. Then*

$$Q_\varphi = conv\left(\bigcup_{k=0}^{\nu} S_k\right) + \{(y, z) \in \mathbb{R}_+ \times \mathbb{R}^n : z = \mathbf{0}\}. \tag{6}$$

Proof. Observe that

$$C_\varphi = \left\{(\gamma, \alpha, \beta) \in \mathbb{R}_+ \times \mathbb{R}^n \times \mathbb{R} : \gamma y + \sum_{j=1}^{n} \alpha_j z_j \geq \beta, \ \forall \, (y, z) \in \bigcup_{k=0}^{\nu} S_k\right\}.$$

It is therefore clear that

$$S := conv\left(\bigcup_{k=0}^{\nu} S_k\right) + \{(y, z) \in \mathbb{R}_+ \times \mathbb{R}^n : z = \mathbf{0}\} \subseteq Q_\varphi.$$

Suppose, for a contradiction, there exists a point $(y^*, z^*) \in Q_\varphi - S$. Then there is a valid inequality $\gamma^* y + \sum_{j=1}^{n} \alpha_j^* z_j \geq \beta^*$ for S that is violated by (y^*, z^*). However, $(\gamma^*, \alpha^*, \beta^*) \in C_\varphi$, implying that

$$\gamma^* y^* + \sum_{j=1}^{n} \alpha_j^* z_j^* \geq \beta^*,$$

as $(y^*, z^*) \in Q_\varphi$, a contradiction. Therefore, $S = Q_\varphi$, as required. □ □

Given the results in Theorem 4, we can now apply Balas's theory of disjunctive programming [2, 3] to (6), enabling us to get the following extended formulation for Q_φ. Recall that $R_k = \{z \in [0, 1]^n : A^k z \leq b^k\}$ for $0 \leq k \leq \nu$.

Corollary 5. $Q_\varphi = proj_{y,z} EQ_\varphi$ *for*

$$EQ_\varphi := \left\{(y, z, \lambda, \omega) \in \mathbb{R}_+ \times [0, 1]^n \times \mathbb{R}_+^{\nu+1} \times \mathbb{R}_+^{(\nu+1)n} : (7) - (11)\right\}$$

where

$$\sum_{k=0}^{\nu} \lambda_k = 1 \tag{7}$$

$$\sum_{k=0}^{\nu} \omega^k = z \tag{8}$$

$$\omega^k \leq \lambda_k \mathbf{1}, \quad \forall 0 \leq k \leq \nu \tag{9}$$

$$A^k \omega^k \leq b^k \lambda_k, \quad \forall 0 \leq k \leq \nu \tag{10}$$

$$y \geq \sum_{k=0}^{\nu} h_{k+1} \lambda_k. \tag{11}$$

Moreover, we can actually get an extended formulation for C_φ as well, according to the following proposition.

Proposition 6. $C_\varphi = proj_{\gamma,\alpha,\beta} EC_\varphi$ for

$$EC_\varphi := \left\{ (\gamma, \alpha, \beta, \sigma, \rho) \in \mathbb{R}_+ \times \mathbb{R}^n \times \mathbb{R} \times \mathbb{R}_-^{\sum_{k=0}^\nu n_k} \times \mathbb{R}_-^{(\nu+1)n} : (12); (13) \right\}$$

where

$$\gamma h_{k+1} + \sum_{j=1}^k \alpha_j + (b^k)^\top \sigma^k + (\rho^k)^\top \mathbf{1} \geq \beta, \ \forall\, 0 \leq k \leq \nu \tag{12}$$

$$(A^k)^\top \sigma^k + \rho^k \leq \alpha^k, \ \forall\, 0 \leq k \leq \nu \tag{13}$$

and where, for each $0 \leq k \leq \nu$, $\alpha^k = (0, 0, \ldots, 0, \alpha_{k+1}, \alpha_{k+2}, \ldots, \alpha_n) \in \mathbb{R}^n$.

Proof. For $0 \leq k \leq \nu$ and $\alpha \in \mathbb{R}^n$, we know that

$$\begin{aligned}
\varphi(k, \alpha) = \quad &\min \quad \sum_{j=k+1}^n \alpha_j z_j \\
&\text{s.t.} \quad A^k z \leq b^k \\
&\qquad\ z \leq \mathbf{1} \\
&\qquad\ z \geq \mathbf{0}.
\end{aligned}$$

which, by strong LP duality, is equal to

$$\begin{aligned}
\varphi(k, \alpha) = \quad &\max \quad (b^k)^\top \sigma^k + (\rho^k)^\top \mathbf{1} \\
&\text{s.t.} \quad (A^k)^\top \sigma^k + \rho^k \leq \alpha^k \\
&\qquad\ \sigma^k, \rho^k \leq \mathbf{0}.
\end{aligned}$$

Note now that weak LP duality implies $C_\varphi \supseteq proj_{\gamma,\alpha,\beta} EC_\varphi$, and by strong LP duality, we have $C_\varphi \subseteq proj_{\gamma,\alpha,\beta} EC_\varphi$. Therefore, $C_\varphi = proj_{\gamma,\alpha,\beta} EC_\varphi$. $\qquad\qquad\square\qquad\qquad\qquad\qquad\square$

We note that these results generalize the previous results of [14, 12] where the extended formulations were considered by using the LP relaxation of the cardinality constrained case, which in that case coincides with the convex hull of the knapsack polytope.

3.2 On the strength of the relaxations

Here we consider what can be said about the strength of the proposed relaxations, given knowledge about the strength of the LP relaxations of the knapsack polytope. Our first theorem shows that if the relaxation is exact for a given subset of variables, then all the inequalities that are allowed to use negative coefficients for the same subset of variables are guaranteed to be considered in the relaxation.

To do so, let us define for any subset $J \subseteq [n]$ the sets

$$R_k^J = R_k \cap \{ z \in \mathbb{R}^n : z_j = 0, \forall j \notin J : j > k \}$$

$$P_k^J = P_k \cap \{ z \in \mathbb{R}^n : z_j = 0, \forall j \notin J : j > k \},$$

that is, the sets R_k and P_k when we restrict ourselves only to variables in J. We now want to consider what happens when the relaxations R_k^J are exact.

Theorem 7. *Suppose, for a given subset $J \subseteq [n]$, we have that*

$$\min\left\{ \sum_{j=k+1}^{n} \alpha_j z_j : z \in R_k^J \right\} = \min\left\{ \sum_{j=k+1}^{n} \alpha_j z_j : z \in P_k^J \right\}, \quad \forall \alpha \in \mathbb{R}^n.$$

Take $(y^*, z^*) \in Q_\varphi$. Then $\gamma y^* + \sum_{j \in [n]} \alpha_j z_j^* \geq \beta$, for all $(\gamma, \alpha, \beta) \in C$ such that $\{j \in [n] : \alpha_j < 0\} \subseteq J$.

Proof. Choose $(\gamma, \alpha, \beta) \in C$ such that $\{j \in [n] : \alpha_j < 0\} \subseteq J$. Observe that, in this case, the argument of $\phi(k, \alpha)$ will be in R_k^J since for all j with $\alpha_j \geq 0$ we can just set $z_j = 0$ and either improve or not change the value of $\phi(k, \alpha)$. Similarly, the argument of $\varphi(k, \alpha)$ will be in P_k^J. Therefore, we have that $\phi(k, \alpha) = \varphi(k, \alpha)$ for all $0 \leq k \leq \nu$, and so $(\gamma, \alpha, \beta) \in C_\varphi$. As a consequence, since $(y^*, z^*) \in Q_\varphi$, we have that $\gamma y^* + \sum_{j \in [n]} \alpha_j z_j^* \geq \beta$, as desired. $\qquad\qquad\qquad\square$ $\qquad\qquad\square$

While this theorem may seem strange at first, what it is saying is that as long as we can describe completely the convex hull of the knapsack polytope restricted to a set J, then we are guaranteed that all the points in Q_φ satisfy all the inequalities where the only negative coefficients allowed are in J. For instance, when we are in the cardinality constrained case, we have that $J = [n]$ and so this says that all inequalities of $conv(Q)$ must be satisfied by Q_φ. As another special case, for instance, one can derive the following corollary.

Corollary 8. *Suppose that*

$$R_k \subseteq \left\{ z \in [0,1]^n : \begin{array}{l} \sum_{j=k+1}^{n} \lfloor d a_j \rfloor z_j \leq \left\lfloor d\left(p - \sum_{j=1}^{k} a_j\right)\right\rfloor, \\ \forall d \in \left\{ \frac{1}{a_j} : j \in [n] \right\} \end{array} \right\}.$$

Take $(y^*, z^*) \in Q_\varphi$. Then, for each $(\gamma, \alpha, \beta) \in C$ with the property that $a_i = a_j$ for all $\alpha_i, \alpha_j < 0$, we have

$$\gamma y^* + \sum_{j \in [n]} \alpha_j z_j^* \geq \beta.$$

In particular, if $a_1 = a_2 = \cdots = a_n$ then $Q_\varphi = conv(Q)$.

Next we study the case where φ is guaranteed to approximate ϕ. More precisely, suppose there exists an $\delta \in (0,1)$ such that for every $0 \leq k \leq \nu$,

$$(\diamond) \quad \max\left\{ w^\top z : z \in P_k \right\} \geq (1-\delta)\max\left\{ w^\top z : z \in R_k \right\} \text{ for all } w \in \mathbb{R}^n.$$

That is, we have that for each $0 \leq k \leq \nu$ and $\alpha \in \mathbb{R}^n$,

$$(\diamond') \quad (1-\delta)\varphi(k,\alpha) \geq \phi(k,\alpha) \geq \varphi(k,\alpha).$$

(Recall $\phi(k,\alpha) \leq 0$.) We now show that such approximation factor can be translated into the inequalities in the following way. For any $\alpha \in \mathbb{R}^n$, let $S^-(\alpha) := \sum \left(\alpha_j : 1 \leq j \leq n, \alpha_j < 0 \right)$. Then, the next lemma shows that we are guaranteed that all inequalities for $conv(Q)$ are not violated by much if we consider Q_φ.

Proposition 9. *If $(\gamma, \alpha, \beta) \in C$ and $R_0, R_1, \ldots, R_\nu$ satisfy $(\diamond)$, then $(\gamma, \alpha, \beta + \delta S^-(\alpha)) \in C_\varphi$.*

Proof. Take $(\gamma, \alpha, \beta) \in C$. By $(\diamond')$, for each $0 \leq k \leq \nu$,

$$\gamma h_{k+1} + \sum_{j \leq k} \alpha_j + \varphi(k, \alpha) \geq \gamma h_{k+1} + \sum_{j \leq k} \alpha_j + \phi(k, \alpha) + \delta\varphi(k, \alpha) \geq \beta + \delta\varphi(k, \alpha).$$

On the other hand,

$$\delta\varphi(k, \alpha) = \delta \min \left\{ \sum_{j > k} \alpha_j z_j^* : z^* \in R_k \right\} \geq \delta S^-(\alpha),$$

so

$$\gamma h_{k+1} + \sum_{j \leq k} \alpha_j + \varphi(k, \alpha) \geq \beta + \delta S^-(\alpha).$$

Since this is true for all $0 \leq k \leq \nu$, it follows that $(\gamma, \alpha, \beta + \delta S^-(\alpha)) \in C_\varphi$. $\quad\square\quad\quad\quad\square$

One application of Proposition 9, for instance, is to consider Bienstock's approximate formulation for the knapsack polytope. As shown in Bienstock [8], for every $0 \leq k \leq \nu$ and $\delta > 0$, there exists a polyhedron R_k^δ of dimension $O\big(\delta^{-1} n^{1+\lceil 1/\delta \rceil}\big)$ that is described by $O\big(\delta^{-1} n^{2+\lceil 1/\delta \rceil}\big)$ constraints whose projection onto z satisfies the desired property $(\diamond)$.

4 Facet-defining inequalities for $conv(Q)$

The results in the previous section allow us to study several different relaxations of $conv(Q)$ by studying relaxations of the knapsack polytope.

In this section, we focus on developing new classes of facet-defining inequalities for $conv(Q)$ that do not arise from facet-defining inequalities for the knapsack set P_0. By Theorem 1 and Proposition 2, such inequalities must have the form $y + \sum_{j \in [n]} \alpha_j z_j \geq \beta$.

We begin this section by giving an overview of the known classes of facet-defining inequalities for $conv(Q)$. We then describe a property of facet-defining inequalities for $conv(Q)$ that generalizes the results for the previously known facet-defining inequalities. Finally, we will introduce a new explicit class of facet-defining inequalities that subsumes all the previously known classes.

The first proposed class of facet-defining inequalities for $conv(Q)$ were the so-called *strengthened star inequalities*.

Theorem 10 ([14])**.** *The strengthened star inequalities*

$$y + \sum_{j=1}^{a} (h_{t_j} - h_{t_{j+1}}) z_{t_j} \geq h_{t_1}, \; \forall \, T = \{t_1, \dots, t_a\} \subseteq \{1, \dots, \nu\} \tag{14}$$

with $t_1 < \dots < t_a$ and $h_{t_{a+1}} := h_{\nu+1}$ are valid for $conv(Q)$. Moreover, (14) is facet-defining for $conv(Q)$ if and only if $h_{t_1} = h_1$. $\quad\square$

As shown in [10, 1, 9], the strengthened star inequalities can be separated in polynomial time and are sufficient to describe the convex hull of

$$\left\{ (y,z) \in \mathbb{R}_+ \times \{0,1\}^n : y + h_i z_i \geq h_i, \ \forall \, i \in [n] \right\}.$$

However, as one may expect, when a knapsack constraint is enforced (to obtain Q), the convex hull becomes much more complex, so the facet-defining inequalities become more difficult to find.

In 2010, Luedtke et al. [14] found a general class of facet-defining inequalities for $conv(Q)$. Subsequently, Küçükyavuz [12] introduced a larger and subsuming class of facet-defining inequalities for $conv(Q)$, called the (T, Π_L) *inequalities*. Here, we only state the latter class.

Theorem 11 ([12]). *Suppose that $a_1 = \cdots = a_n = 1$, p is a positive integer and $m \in [p]$. Suppose further that*

(i) $T := \{t_1, \ldots, t_a\} \subseteq [m]$ *where* $t_1 < \cdots < t_a$,

(ii) $L \subseteq \{m+2, \ldots, n\}$ *is of size* $p - m$ *and* $\Pi_L := (\ell_1, \ldots, \ell_{p-m})$ *is a permutation of the elements of L such that $\ell_j > m + j$, for $j \in [p - m]$.*

Set $t_{a+1} := m + 1$. Let $\Delta_1 := h_{m+1} - h_{m+2}$, and for $2 \leq j \leq p - m$, define

$$\Delta_j := \max\left\{ \Delta_{j-1}, h_{m+1} - h_{m+1+j} - \sum\left(\Delta_i : \ell_i > m+j, i < j \right) \right\}.$$

Then the (T, Π_L) inequality

$$y + \sum_{j=1}^{a}(h_{t_j} - h_{t_{j+1}})z_{t_j} + \sum_{j=1}^{p-m} \Delta_j(1 - z_{\ell_j}) \geq h_{t_1} \tag{15}$$

is valid for $conv(Q)$. Furthermore, (15) is facet-defining inequality for $conv(Q)$ if and only if $h_{t_1} = h_1$. □

Observe that the $(T, \emptyset)$ inequalities are simply the strengthened star inequalities. A common aspect of Theorems 10 and 11 (and the other class of facet-defining inequalities in Luedtke et al. [14]) is the condition on the right-hand-side constant for the inequality to be facet-defining. The following result provides a generalization of those conditions to any facet-defining inequality of $conv(Q)$.

Proposition 12. *If*

$$y + \sum_{j\in[n]} \alpha_j z_j \geq \beta \tag{16}$$

is valid for $conv(Q)$ then $\beta \leq h_1 + \phi(0, \alpha)$. Moreover, if (16) is facet-defining for $conv(Q)$, then $\beta = h_1 + \phi(0, \alpha)$.

Proof. Take $z^* \in P_0$ so that $\sum_{j\in[n]} \alpha_j z_j^* = \phi(0, \alpha)$. As $(h_1, z^*) \in Q$ we have

$$h_1 + \phi(0, \alpha) = h_1 + \sum_{j\in[n]} \alpha_j z_j^* \geq \beta.$$

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Suppose now that (16) is facet-defining. Since $conv(Q)$ is full-dimensional and (16) is a facet-defining inequality different from $z_1 \leq 1$, it follows that there is a point $(y', z') \in Q$ on the facet defined by (16) such that $z'_1 = 0$. Then $y' = y' + h_1 z'_1 \geq h_1$, and since $z' \in P_0$, it follows that

$$\beta = y' + \sum_{j \in [n]} \alpha_j z'_j \geq h_1 + \phi(0, \alpha).$$

Hence, $\beta = h_1 + \phi(0, \alpha)$. $\qquad\qquad\qquad\qquad\qquad\qquad \square \qquad\qquad\qquad\qquad\qquad\qquad \square$

We note that for (14) we have $\phi(0, \alpha) = 0$ and for (15) we have $\phi(0, \alpha) = -\sum_{i=1}^{p-m} \Delta_i$, so the above theorem implies the previously stated results on the right-hand-side constants of the inequalities.

We now introduce a class of valid inequalities for $conv(Q)$ subsuming all the preceding classes of facet-defining inequalities. Let $s_0 := 0$ and for each $m \in [n]$, let $s_m := \sum_{j=1}^{m} a_j$.

Theorem 13. *Choose $m \in \{0, 1, \ldots, \nu\}$ such that $p - s_m \leq n - m - 1$. Let s be an integer such that $p - s_m \leq s \leq n - m - 1$. For $1 \leq j \leq s$, let $\mathfrak{m}(j) := \max\{k : k \in [n], j \geq s_k - s_m\}$. Suppose that*

(V1) $T := \{t_1, \ldots, t_a\} \subseteq \{1, \ldots, m\}$ where $t_1 < \cdots < t_a$,

(V2) $L := \{\ell_1, \ldots, \ell_s\} \subseteq \{m+2, \ldots, n\}$ where $\ell_1, \ldots, \ell_s$ are pairwise distinct,

(V3) $a_j \geq 1$ for all $j \in L$.

Set $t_{a+1} := m + 1$. Choose $\Delta \in \mathbb{R}^L$ such that $0 \leq \Delta_{\ell_1} \leq \Delta_{\ell_2} \leq \cdots \leq \Delta_{\ell_s}$ and

$$\sum \left(\Delta_{\ell_i} : \ell_i > \mathfrak{m}(j), i \leq j \right) - \sum \left(\Delta_{\ell_i} : \ell_i \leq \mathfrak{m}(j), i > j \right) \geq h_{m+1} - h_{\mathfrak{m}(j)+1}$$

for all $1 \leq j \leq s$. Then

$$y + \sum_{j=1}^{a} (h_{t_j} - h_{t_{j+1}}) z_{t_j} + \sum_{i \in L} \Delta_i (1 - z_i) \geq h_{t_1} \qquad (17)$$

is a valid inequality for $conv(Q)$.

A proof is included in §4.2. Under certain conditions described below, (17) becomes a facet-defining inequality for $conv(Q)$.

Theorem 14. *Choose $m \in \{0, 1, \ldots, \nu\}$ such that $p - s_m \leq n - m - 1$ and $p - s_m$ is an integer. For $1 \leq j \leq p - s_m$, let $\mathfrak{m}(j) := \max\{k : k \in [n], j \geq s_k - s_m\}$. Suppose that*

(F1) $T := \{t_1, \ldots, t_a\} \subseteq \{1, \ldots, m\}$ where $t_1 < \cdots < t_a$ and $h_{t_1} = h_1$,

(F2) $L \subseteq \{m+2, \ldots, n\}$ is of size $p - s_m$ and $(\ell_1, \ldots, \ell_{p-s_m})$ is a permutation of the elements of L such that $\ell_j > \mathfrak{m}(j)$, for all $1 \leq j \leq p - s_m$,

(F3) $a_j = 1$ for all $j \in L$, and $a_i \leq s_m$ for all $i \in [n] - L$.

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Set $t_{a+1} := m + 1$. *Let* $\Delta_{\ell_1} := h_{m+1} - h_{\mathfrak{m}(1)+1}$, *and for* $2 \le j \le p - s_m$, *define*

$$\Delta_{\ell_j} := \max \left\{ \Delta_{\ell_{j-1}}, h_{m+1} - h_{\mathfrak{m}(j)+1} - \sum \left(\Delta_{\ell_i} : \ell_i > \mathfrak{m}(j), i < j \right) \right\}.$$

Then (17) *is a facet-defining inequality for* $conv(Q)$.

A proof can be found in §4.3. The class above coincides with the (T, Π_L) inequalities in the case when $a_1 = \cdots = a_n = 1$ and p is an integer.

Let us explain how this theorem implies that the (T, Π_L) inequalities (15) are facet-defining for $conv(Q)$. Observe that, in the context of Theorem 11, $s_j = j$, for all $0 \le j \le n$. Let $s := p - m$ and note that $\mathfrak{m}(j) = m + j$, for all $1 \le j \le s$. Furthermore, $\ell_j > m + j = \mathfrak{m}(j)$ for all $1 \le j \le s$. Hence, by Theorem 14, the (T, Π_L) inequalities (15) are facet-defining for $conv(Q)$.

Observe that Theorems 13 and 14 can also be applied to any scalar multiple of the knapsack constraint, and this will potentially give us more facet-defining inequalities. That is, one can apply Theorem 14 to $\sum_{j \in [n]} da_j z_j \le dp$, for any arbitrary positive real number d. This is explained in the following example.

Example 1. *Let* $a = (2, 1.5, 2.5, 1, 1, 1, 2, 1, 0.5, 0.5)$ *and*
$h = (809, 405, 202, 100, 60, 40, 30, 25, 23, 20)$ *for* $n = 10$ *and* $p = 9$. *Note that* $\nu = 6$. *For* $m = 6$, $T = \{1, 3, 5\}$ *and* $L = \emptyset$, *the (strengthened star) inequality*

$$y + (809 - 202)z_1 + (202 - 60)z_3 + (60 - 30)z_5 \ge 809$$

is facet-defining for $conv(Q)$.

Next set $m = 3$, $T = \{1, 2, 3\}$ *and* $L = \{\ell_1 = 5, \ell_2 = 6, \ell_3 = 8\}$. *Then* $\mathfrak{m}(1) = 4, \mathfrak{m}(2) = 5, \mathfrak{m}(3) = 6$ *and the hypotheses of Theorem 14 are satisfied. Thus the inequality*

$$y + (809 - 405)z_1 + (405 - 202)z_2 + (202 - 100)z_3$$
$$+ 40(1 - z_5) + 60(1 - z_6) + 70(1 - z_8) \ge 809$$

is facet-defining for $conv(Q)$.

Furthermore, one can replace a *and* p *with the equivalent choice of* $2a$ *and* $2p$. *In this case, set* $m = 5$, $T = \{1, 2, 5\}$ *and* $L = \{\ell_1 = 9, \ell_2 = 10\}$. *Then* $\mathfrak{m}(1) = 5, \mathfrak{m}(2) = 6$ *and the hypotheses of Theorem 14 are once again satisfied. Thus the inequality*

$$y + (809 - 405)z_1 + (405 - 60)z_2 + (60 - 40)z_5 + 0(1 - z_9) + 10(1 - z_{10}) \ge 809$$

is facet-defining for $conv(Q)$.

4.1 Separation of a subset of proposed facet-defining inequalities

Separating over all proposed facet-defining inequalities (17) seems to be hard, the bottleneck being minimizing for a fixed z^* the expression $\sum_{i \in L} \Delta_i (1 - z_i^*)$ over all possible L's, as the choice of L affects the values of Δ. To circumvent this difficulty, Küçükyavuz [12] restricted the choices for L so as to control the values of Δ. With not much more work, we can obtain a similar result for our general setting:

Proposition 15. *There is an exact separation algorithm with running time $O(p^4)$ that separates over the proposed facet-defining inequalities* (17) *for which there is a partition of L into parts F, G where the following hold:*

(1) $\mathfrak{m}(1) < \mathfrak{m}(2) < \cdots < \mathfrak{m}(p - s_m)$,

(2) $F = \{\mathfrak{m}(1) + 1, \ldots, \mathfrak{m}(r) + 1\}$ *for some* $1 \leq r \leq p - s_m$,

(3) $G \subseteq \{\mathfrak{m}(p - s_m) + 1, \mathfrak{m}(p - s_m) + 2, \ldots, n\}$.

The proof of this result is almost identical to that in [12], so we refrain from including it.

4.2   Proof of Theorem 13

We prove Theorem 13 using the following two claims. Let $R := [n] - L$ and define $\alpha \in \mathbb{R}^n$ as follows:

$$\alpha_i := \begin{cases} h_{t_j} - h_{t_{j+1}}, & \text{if } i = t_j \text{ for some } 1 \leq j \leq a; \\ 0, & \text{if } i \in R - T; \\ -\Delta_i, & \text{if } i \in L. \end{cases}$$

Claim 1. *We have*

$$\phi(k, \alpha) \geq \begin{cases} \sum_{j \in L} \alpha_j, & \text{if } 0 \leq k \leq m; \\ \sum (\alpha_{\ell_i} : i > s_k - s_m), & \text{if } m + 1 \leq k \leq \nu. \end{cases}$$

Proof of Claim. Since $\alpha_i < 0$ only for indices $i \in L$, it follows immediately that, for $0 \leq k \leq m$,

$$\phi(k, \alpha) \geq \sum_{j \in L} \alpha_j.$$

Next choose $m + 1 \leq k \leq \nu$ and let $z \in P_k$. Then

$$|\{j \in L : j > k, z_j = 1\}| = \sum_{j \in L, j > k} z_j \leq \sum_{j \in L, j > k} a_j z_j \quad \text{by } (V3)$$

$$\leq p - \sum_{j \leq k} a_j$$

$$= p - s_k$$

$$\leq |L| - s_k + s_m$$

$$= |\{\ell_i \in L : i > s_k - s_m\}|.$$

Therefore, as $\alpha_i \geq 0$ for all $i \in [n] - L$ and $0 \geq \alpha_{\ell_1} \geq \cdots \geq \alpha_{\ell_{p-s_m}}$, it follows

$$\sum_{j > k} \alpha_j z_j \geq \sum_{j \in L, j > k} \alpha_j z_j \geq \sum (\alpha_{\ell_i} : i > s_k - s_m).$$

Since this is true for all $z \in P_k$, we must have

$$\phi(k, \alpha) \geq \sum (\alpha_{\ell_i} : i > s_k - s_m),$$

as claimed.   ◊

Claim 2. $\left(1, \alpha, h_{t_1} + \sum_{i \in L} \alpha_i\right) \in C$.

Proof of Claim. Let $0 \le k \le \nu$. If $0 \le k \le m$, choose $j \in \{0, 1, \ldots, a\}$ so that $t_j < k+1 \le t_{j+1}$, where $t_0 := 0$. In this case, $\phi(k, \alpha) \ge \sum_{i \in L} \alpha_i$ by Claim 1, so

$$h_{k+1} + \sum_{i=1}^{k} \alpha_i + \phi(k, \alpha) \ge h_{t_{j+1}} + \sum_{i=1}^{j}(h_{t_i} - h_{t_{i+1}}) + \sum_{i \in L} \alpha_i$$

$$= h_{t_1} + \sum_{i \in L} \alpha_i.$$

Otherwise we have $m + 1 \le k \le \nu$. Then

$$\sum_{i \in R, i \le k} \alpha_i = \sum_{i \in T} \alpha_i = h_{t_1} - h_{m+1}.$$

By Claim 1,

$$\phi(k, \alpha) \ge \sum \left(\alpha_{\ell_i} : i > s_k - s_m\right).$$

Let $j := s_k - s_m$. Note that $\mathfrak{m}(j) = \max\{r : j \ge s_r - s_m\} = k$ as $a_i > 0$ for all $i \in [n]$. Then

$$h_{k+1} + \sum_{i=1}^{k} \alpha_i + \phi(k, \alpha)$$

$$= h_{k+1} + \sum_{i \in R, i \le k} \alpha_i + \sum_{i \in L, i \le k} \alpha_i + \phi(k, \alpha)$$

$$\ge h_{k+1} + h_{t_1} - h_{m+1} + \sum \left(\alpha_{\ell_i} : \ell_i \le k\right) + \sum \left(\alpha_{\ell_i} : i > j\right)$$

$$= h_{t_1} + h_{k+1} - h_{m+1} + \sum_{i \in L} \alpha_i - \sum \left(\alpha_{\ell_i} : \ell_i > k, i \le j\right) + \sum \left(\alpha_{\ell_i} : \ell_i \le k, i > j\right)$$

$$= h_{t_1} + h_{k+1} - h_{m+1} + \sum_{i \in L} \alpha_i + \sum \left(\Delta_{\ell_i} : \ell_i > k, i \le j\right) - \sum \left(\Delta_{\ell_i} : \ell_i \le k, i > j\right)$$

$$\ge h_{t_1} + h_{k+1} - h_{m+1} + \sum_{i \in L} \alpha_i + h_{m+1} - h_{k+1}$$

$$= h_{t_1} + \sum_{i \in L} \alpha_i.$$

As a result, $\left(1, \alpha, h_{t_1} + \sum_{i \in L} \alpha_i\right) \in C$, proving Claim 2. $\diamond$

Observe that Claim 2 implies that (17) is a valid inequality for $conv(Q)$, as (17) is equivalent to

$$y + \sum_{j \in [n]} \alpha_j z_j \ge h_{t_1} + \sum_{i \in L} \alpha_i. \tag{18}$$

$\square$

4.3 Proof of Theorem 14

As in the proof of Theorem 13, let $R := [n] - L$ and define $\alpha \in \mathbb{R}^n$ as follows:

$$\alpha_i := \begin{cases} h_{t_j} - h_{t_{j+1}}, & \text{if } i = t_j \text{ for some } 1 \le j \le a; \\ 0, & \text{if } i \in R - T; \\ -\Delta_i, & \text{if } i \in L. \end{cases}$$

We first show that (17) is a valid inequality for $conv(Q)$. Observe first that (V1)-(V3) are satisfied. It is clear by definition that $0 \le \Delta_{\ell_1} \le \cdots \le \Delta_{\ell_s}$. Moreover, for each $1 \le j \le p - s_m$,

$$h_{m+1} - h_{\mathfrak{m}(j)+1} \le \Delta_{\ell_j} + \sum \left(\Delta_{\ell_i} : \ell_i > \mathfrak{m}(j), i < j \right)$$

$$= \sum \left(\Delta_{\ell_i} : \ell_i > \mathfrak{m}(j), i \le j \right) \qquad (\dagger)$$

$$= \sum \left(\Delta_{\ell_i} : \ell_i > \mathfrak{m}(j), i \le j \right) - \sum \left(\Delta_{\ell_i} : \ell_i \le \mathfrak{m}(j), i > j \right). \; (\ddagger)$$

Above, $(\dagger)$ holds since $\ell_j > \mathfrak{m}(j)$, and $(\ddagger)$ holds since $\ell_i > \mathfrak{m}(i) \ge \mathfrak{m}(j)$ for all $i > j$. As a result, Theorem 13 applies and implies that (17) is indeed a valid inequality for $conv(Q)$.

We will next find $n + 1$ affinely independent points in Q that satisfy (17), or equivalently (18), at equality:

1. For each $k := t_j \in T$, let $y^k := h_k$ and define $z^k \in \{0, 1\}^n$ as follows:

$$z_i^k := \begin{cases} 1, & \text{if } i < k \text{ or } i \in L; \\ 0, & \text{otherwise.} \end{cases}$$

We have $(y^k, z^k) \in Q$, because $z_i^k = 1$ for $i = 1, \ldots, k - 1$, and

$$\sum_{i=1}^n a_i z_i^k = \sum_{i<k} a_i z_i^k + \sum_{i \in L} z_i^k = s_{k-1} + |L| = s_{k-1} + p - s_m \le p.$$

Moreover, (y^k, z^k) satisfies (18) at equality:

$$y^k + \sum_{i=1}^a (h_{t_i} - h_{t_{i+1}}) z_{t_i}^k + \sum_{i \in L} \alpha_i z_i^k = h_{t_j} + \sum_{i=1}^{j-1} (h_{t_i} - h_{t_{i+1}}) + \sum_{i \in L} \alpha_i$$

$$= h_{t_j} + h_{t_1} - h_{t_j} + \sum_{i \in L} \alpha_i$$

$$= h_{t_1} + \sum_{i \in L} \alpha_i.$$

2. For each $k := \ell_j \in L$, let $\mathfrak{f}(j) := \min\{f \in [j] : \Delta_{\ell_j} = \Delta_{\ell_f}\}$. By definition we must have, for each $\ell_j \in L$,

$$\Delta_{\ell_j} = \Delta_{\ell_{\mathfrak{f}(j)}} = h_{m+1} - h_{\mathfrak{m}(\mathfrak{f}(j))+1} - \sum \left(\Delta_{\ell_i} : \ell_i > \mathfrak{m}(\mathfrak{f}(j)), i < \mathfrak{f}(j) \right),$$

and so

$$\alpha_{\ell_j} + \sum \left(\alpha_{\ell_i} : \ell_i > \mathfrak{m}(\mathfrak{f}(j)), i < \mathfrak{f}(j) \right) = h_{\mathfrak{m}(\mathfrak{f}(j))+1} - h_{m+1}. \; (\star)$$

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Now let $y^k := h_{\mathfrak{m}(\mathfrak{f}(j))+1}$ and define $z^k \in \{0,1\}^n$ as follows:

$$z_t^k := \begin{cases} 0, & \text{if } t = \ell_i > \mathfrak{m}(\mathfrak{f}(j)) \text{ and } i < \mathfrak{f}(j), \\ & \quad \text{or } t = \ell_j, \\ & \quad \text{or } t \in R \text{ and } t > \mathfrak{m}(\mathfrak{f}(j)); \\ 1, & \text{otherwise.} \end{cases}$$

Observe that $\ell_j > \mathfrak{m}(j) \geq \mathfrak{m}(\mathfrak{f}(j))$. We have $(y^k, z^k) \in Q$, because $z_i^k = 1$ for $i = 1, \ldots, \mathfrak{m}(\mathfrak{f}(j))$, and

$$\sum_{i=1}^{n} a_i z_i^k = \sum_{i \leq \mathfrak{m}(\mathfrak{f}(j))} a_i + |\{\ell_i \in L : \ell_i > \mathfrak{m}(\mathfrak{f}(j)), i \geq \mathfrak{f}(j), i \neq j\}|$$

$$= s_{\mathfrak{m}(\mathfrak{f}(j))} + |\{\ell_i \in L : \ell_i > \mathfrak{m}(\mathfrak{f}(j)), i \geq \mathfrak{f}(j)\}| - 1$$

$$= s_{\mathfrak{m}(\mathfrak{f}(j))} + |\{\ell_i \in L : i \geq \mathfrak{f}(j)\}| - 1 \quad (\star\star)$$

$$= s_{\mathfrak{m}(\mathfrak{f}(j))} + |L| - \mathfrak{f}(j)$$

$$= s_{\mathfrak{m}(\mathfrak{f}(j))} + p - s_m - \mathfrak{f}(j)$$

$$\leq p \quad \text{by the definition of } \mathfrak{m}.$$

Above, $(\star\star)$ holds because $i \geq \mathfrak{f}(j)$ implies that $\ell_i > \mathfrak{m}(i) \geq \mathfrak{m}(\mathfrak{f}(j))$. Moreover, (y^k, z^k) satisfies (18) at equality:

$$y^k + \sum_{i=1}^{a}(h_{t_i} - h_{t_{i+1}})z_{t_i}^k + \sum_{i \in L} \alpha_i z_i^k$$

$$= h_{\mathfrak{m}(\mathfrak{f}(j))+1} + h_{t_1} - h_{m+1} + \sum_{i \in L} \alpha_i - \alpha_{\ell_j} - \sum (\alpha_{\ell_i} : \ell_i > \mathfrak{m}(\mathfrak{f}(j)), i < \mathfrak{f}(j))$$

$$= h_{\mathfrak{m}(\mathfrak{f}(j))+1} + h_{t_1} - h_{m+1} + \sum_{i \in L} \alpha_i + h_{m+1} - h_{\mathfrak{m}(\mathfrak{f}(j))+1} \quad \text{by } (\star)$$

$$= h_{t_1} + \sum_{i \in L} \alpha_i.$$

3. For all $k \in R - T$, let $y^k := h_1$ and define $z^k \in \{0,1\}^n$ as follows:

$$z_i^k := \begin{cases} 0, & \text{if } i \in R - \{k\}; \\ 1, & \text{otherwise.} \end{cases}$$

Notice $(y^k, z^k) \in Q$, because

$$\sum_{i=1}^{n} a_i z_i^k = a_k + |L| = a_k + p - s_m \leq p$$

by (F3). Moreover,

$$y^k + \sum_{i=1}^{a}(h_{t_i} - h_{t_{i+1}})z_{t_i}^k + \sum_{i \in L} \alpha_i z_i^k = h_1 + \sum_{i \in L} \alpha_i = h_{t_1} + \sum_{i \in L} \alpha_i.$$

4. Lastly, let $y^0 := h_{m+1}$ and define $z^0 \in \{0,1\}^n$ as follows:

$$z_i^0 := \begin{cases} 1, & \text{if } i < m+1 \text{ or } i \in L; \\ 0, & \text{otherwise.} \end{cases}$$

Then $(y^0, z^0) \in Q$ because $z_i^0 = 1$ for all $i < m+1$, and

$$\sum_{i=1}^{n} a_i z_i^k = \sum_{i<m+1} a_i z_i^k + \sum_{i \in L} z_i^k = s_m + |L| = s_m + p - s_m = p.$$

Moreover,

$$y^0 + \sum_{i=1}^{a}(h_{t_i} - h_{t_{i+1}})z_{t_i}^0 + \sum_{i \in L} \alpha_i z_i^0 = h_{m+1} + \sum_{i=1}^{a}(h_{t_i} - h_{t_{i+1}}) + \sum_{i \in L} \alpha_i$$

$$= h_{m+1} + h_{t_1} - h_{m+1} + \sum_{i \in L} \alpha_i$$

$$= h_{t_1} + \sum_{i \in L} \alpha_i.$$

Hence, the face defined by (17) contains $(y^0, z^0), (y^1, z^1), \ldots, (y^n, z^n)$, which are, by a routine argument, affinely independent points in Q. As a result, (17) is a facet-defining inequality for $conv(Q)$. $\qquad \square$

5 A compact extended formulation for (CMIP)

So far we have only focused on the single-constraint chance-constrained problem. We now briefly turn to the general problem. In this section, we propose an extended formulation for the set given by

$$y = Bx \tag{19}$$

$$x \in X \tag{20}$$

$$y + \xi^j z_j \geq \xi^j, \ \forall \, j \in [n] \tag{21}$$

$$\sum_{j=1}^{n} a_j z_j \leq p \tag{22}$$

$$z \in \{0,1\}^n. \tag{23}$$

Let $Q_\kappa := \mathcal{D}_\kappa$ and $h_{j,\kappa} := \xi_\kappa^j$ for all $j \in [n]$ and $\kappa \in [d]$. Let $(1_\kappa, 2_\kappa, \ldots, n_\kappa)$ be a permutation of $[n]$ such that $h_{1_\kappa,\kappa} \geq h_{2_\kappa,\kappa} \geq \ldots \geq h_{n_\kappa,\kappa}$, for all $\kappa \in [d]$. Let $\nu_\kappa := \max\{t : \sum_{j \leq t} a_{j_\kappa} \leq p\}$, for all $\kappa \in [d]$.

Theorem 16. *The formulation*

$$\sum_{j=1}^{\nu_\kappa+1} \lambda_{j_\kappa,\kappa} = 1, \; \forall \, \kappa \in [d] \tag{24}$$

$$y_\kappa - \sum_{j=1}^{\nu_\kappa+1} h_{j_\kappa,\kappa} \lambda_{j_\kappa,\kappa} \geq 0, \; \forall \, \kappa \in [d] \tag{25}$$

$$z_i - \sum_{j=1}^{\nu_\kappa+1} \omega_{j_\kappa,\kappa}^i = 0, \; \forall \, \kappa \in [d], i \in [n] \tag{26}$$

$$\left(p - \sum_{i=1}^{j-1} a_{i_\kappa} \right) \lambda_{j_\kappa,\kappa} - \sum_{i=j}^{n} a_{i_\kappa} \omega_{j_\kappa,\kappa}^{i_\kappa} \geq 0, \; \forall \, \kappa \in [d], j \in [\nu_\kappa + 1] \tag{27}$$

$$\lambda_{j,\kappa} \geq \omega_{j,\kappa}^i \geq 0, \; \forall \, \kappa \in [d], j \in [\nu_\kappa + 1], i \in [n] \tag{28}$$

$$\omega_{j_\kappa,\kappa}^{i_\kappa} \geq \lambda_{j_\kappa,\kappa}, \; \forall \, \kappa \in [d], j \in [\nu_\kappa + 1], i \in [j-1] \tag{29}$$

$$\lambda_{j,\kappa} \geq 0, \; \forall \, \kappa \in [d], j \in [\nu_\kappa + 1] \tag{30}$$

$$\sum_{j=1}^{n} a_j z_j \leq p \tag{31}$$

$$y = Bx \tag{32}$$

$$x \in X \tag{33}$$

$$\lambda \in \{0,1\}^{d+\sum_{\kappa=1}^{d} \nu_\kappa} \tag{34}$$

is an extended formulation for the set given by (19)-(23). *The continuous relaxation of the extended formulation defined by* (24)-(33) *is at least as strong as the continuous relaxation of the formulation defined by* (19)-(22).

We would like to point out that the extended formulation given above is almost the same as the one given in [12], Theorem 8 except that our formulation replaces their (extended knapsack cover) constraint ([12]:32) by constraint (27). It is not difficult to see that neither of the constraints ([12]:32) and (27) dominate the other. However, the proof of Theorem 16 is precisely the same as the proof of that theorem, so we omit it.

6 Computational experiments

One of the advantages of our results is that they allow us to establish benchmarks to identify which are good and useful relaxations to use in the context of chance constrained problems. To illustrate this, we carried out computational experiments to establish how well does one relaxation perform compared with another.

The setup that we chose to set the benchmarks is the same one proposed in [12], that is, the probabilistic lot-sizing problem (originally described in [7]). In this setting, we have that the right-hand-sides ξ_κ^j represent cumulative demands in time period $\kappa = 1, \ldots, d$ under scenario $j = 1, \ldots, n$ and the probabilistic constraint is used to ensure that the probability of shortage of products (or equivalently of not being able to fully satisfy demands) is relatively low.

We generated random instances by considering that all the data is integer and is generated randomly and independently of each other according to the following uniform distributions below:

- the demands in any given time period and in any given scenario are Uniform(1,100),

- the variable production costs are Uniform(1,10),

- the setup costs are Uniform(1,1000)

- the holding costs are Uniform(1,5)

- the coeefficients a_j of the knapsack constraint of Q are Uniform(1,100)

  (the interpretation of these coefficients a_j is that they are a rescaling of the probabilities π_j, which we consider as $\frac{a_j}{\sum\limits_{l=1}^{n} a_l}$).

We generated 20 such random instances for each given combination of values of n (number of scenarios), d (number of time periods) and $\epsilon \in [0, 1]$ (value used to determine the right-hand-side p in the knapsack constraint of Q).

Value p is determined as follows: given a set of knapsack coefficients $\{a_j\}_{j=1}^{n}$ we let $p = \left\lfloor \epsilon \left(\sum\limits_{j=1}^{n} a_j \right) \right\rfloor$. That way the knapsack constraint $\sum\limits_{j=1}^{n} a_j z_j \leq p$ is just a rescaling of the constraint $\sum\limits_{j=1}^{n} \pi_j z_j \leq \epsilon$, and we deal only with integer parameters.

The algorithm that was used to compare the relaxations is described in Algorithm 1 below. In it, some computational choices are implicitly described, such as a minimum violation tolerance in order to declare a cut as violated, the cut normalization and a maximum number of iterations that we allow for the cuts to not improve the bound (to avoid potential issues like cycling).

The separation LP (35) presented in Algorithm 1 is used to separate all inequalities that can be generated using Proposition 6, and that it is used to try to separate the current LP solution from $\mathcal{D}_\kappa$ for a certain $\kappa \in [d]$.

We also note that the algorithm depends on input parameter $\mathcal{K}$ which defines from which of the sets $\mathcal{D}_\kappa$ will the algorithm try to separate the current solution. Besides the choice of $\mathcal{K}$, the algorithm implicitly depends on the choice of knapsack relaxation R_k chosen (in the definition of EC_φ). Therefore we use Algorithm 1 to benchmark different choices of R_k as follows: after completion of Algorithm 1 we compute the final lower bound obtained (and thus the gap) for a particular choice of R_k and use this value to evaluate how good this choice of relaxation is. We explicitly choose to ignore the time, because we do not claim that this is the approach that should be used if you are trying to solve such instances faster. Instead, what our approach serves is to try to identify what is the potential of each choice of relaxation in terms of lower bound improvement. With such results in hand, one can be guided as to where to look for better cuts, knowing that the reason certain classes of cuts are or are not successful is because of the relaxation they are derived from and not because of the success/failure of some separation heuristic.

Algorithm 1: Benchmarking algorithm

Data: Formulation (CMIP) for our problem and a subset $\mathcal{K} \subseteq [d]$

repeat

    Solve the LP relaxation of (CMIP) with any additional cuts that have been found so far.

    Let (x^*, y^*, z^*) be the optimal solution to this relaxation.

    for *every* $\kappa \in \mathcal{K}$ **do**

        Solve the following separation LP:

$$
\begin{aligned}
\min \quad & \gamma y_\kappa^* + \sum_{j \in [n]} \alpha_j z_j^* - \beta \\
\text{s.t.} \quad & (\gamma, \alpha, \beta) \in proj_{\gamma,\alpha,\beta} EC_\varphi \\
& \|(\gamma, \alpha, \beta)\|_\infty \leq 1.
\end{aligned}
\tag{35}
$$

        if *The optimal value is* < -0.001 **then**

            (γ, α, β) defines a cut $\gamma y_\kappa^* + \sum_{j \in [n]} \alpha_j z_j^* \geq \beta$ separating (x^*, y^*, z^*) from $\mathcal{D}_\kappa$ as defined in

            (1) (and thus from the set of feasible solutions to (CMIP)).

            Add such cut to the LP relaxation of (CMIP).

        end

    end

    if *The LP bound did not improve in the last 10 iterations* **then**

        Stop. Finish cut generation.

    end

until *No cuts are found*;

We report the gap closed by the end of the algorithm, calculated as follows: given the optimal (or best known) integer solution value z_{IP} for a given problem, the value z_{LP} of the relaxation of (CMIP) without any cuts and the value z_{final} obtained at completion of Algorithm 1, the gap closed is defined as $\frac{z_{final} - z_{LP}}{z_{IP} - z_{LP}}$.

We benchmarked different relaxations using the same separation procedures. The results are summarized in Table 1. All computations are done on an Intel Core 2 Duo with 2 CPUs of 3.06GHz with 16 GB RAM, using CPLEX 12.4. Below we describe what the columns of Table 1 represent.

Choice of $\mathcal{K}$

Typically, when trying to add cuts to strengthen an LP relaxation, the natural choice for $\mathcal{K}$ is $\mathcal{K} = [d]$. However, for our experiments, we also chose to generate cuts from a single relaxation $\mathcal{D}_\kappa$, that is, for a set $\mathcal{K}$ such that $|\mathcal{K}| = 1$. The reasoning for that is to try to isolate the effect of cuts for $\mathcal{D}_\kappa$ for a single κ from the combined effect that several cuts can have in closing the gap (much like in the general MIP case where cuts generated from different constraints may have a combined effect that may mitigate or boost the advantages that each cut has individually).

We thus tried four different choices for $\mathcal{K}$: $\mathcal{K} = \{1\}$, $\mathcal{K} = \{d/2\}$, $\mathcal{K} = \{d\}$, and $\mathcal{K} = [d]$. These are the values represented in the first column of Table 1 by $\mathcal{K} = 1, d/2, d$ and *All*, respectively.

Columns 2-4

The next three columns in Table 1 represent the data parameters of the instance as described above.

Choice of knapsack relaxation

The next six columns represent the average gaps closed over the 20 instances generated obtained by each choice of relaxation. Column R0 represents the results using only CPLEX generated cuts (i.e. not using Algorithm 1). Columns R1-R5 use only cuts generated using (35) (no CPLEX cuts) for different choices of polyhedral relaxations R_k of the knapsack.

Column R1 presents results defining R_k as the LP relaxation of the knapsack P_k, R2 uses the relaxation proposed in [12] (obtained by using a single extended cover inequality) and R3 combines both relaxations. Note that, even though in principle the results in R3 should dominate R1 and R2, this does not always happen since we abort the cut generation process if there has not been a bound improvement for some prespecified number of iterations.

One can argue that these bounds will probably not be too good since the knapsack relaxations used are not too strong. Therefore, we used strengthened formulations for the knapsack relaxation R_k in bounds R4 and R5. The bound presented in column R4 uses the LP relaxation of the knapsack combined with the Chvátal-Gomory cuts of the form $\sum_{j=1}^{n} \lfloor \frac{a_j}{a_l} \rfloor z_j \leq \lfloor \frac{p}{a_l} \rfloor$, for all $l = 1, \ldots, n$.

Column R5 represents the bound obtained by using the LP relaxation of the knapsack strengthened with cover inequalities. The way we obtained cover inequalities for a relaxation R_k was by using CPLEX's cut generator in the following manner. We set up a completely separate Integer Program whose feasible region is purely $\left\{ z \in \{0,1\}^n : \sum_{j=1}^{n} a_j z_j \leq p \, ; z_1 = \ldots = z_k = 1 \right\}$. We then repeat the following procedure 100 times:

1. Generate random objective function coefficients for each z_j variable with the coefficients being drawn from Uniform(1,100).

2. Turn off all CPLEX preprocessing, heuristics and also all cuts except covers and cliques, which are left on in aggressive mode.

3. Let CPLEX solve only the root node of the Branch-and-cut tree and collect all cuts generated by CPLEX.

4. Add all these cuts to a cut pool (discarding repeated cuts) keeping up to a maximum of 500 cuts in this pool.

The relaxation R_k will then consist of the original LP relaxation of the knapsack strengthened with any cut in the cut pool after this procedure.

We chose to generate cover cuts in this way to make use of CPLEX's existing heuristics for separating such cuts, which are likely much better than anything we could implement for the purpose of this test. Once more we point out that this is a particularly very time consuming process and is not recommended for efficiency.

Finally, column SS represents the gap closed by the strengthened star inequalities [14] (note that the results in columns $R_0, \ldots, R_5$ did not include any strengthened star inequalities).

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$\mathcal{K}$	d	n	ϵ (%)	Gap closed(%)						
				R0	R1	R2	R3	R4	R5	SS
1	10	100	5	48.34	18.15	18.15	18.15	18.15	18.15	4.91
1	10	100	10	46.39	14.29	14.30	14.29	14.29	14.30	5.13
1	10	100	20	41.44	11.78	11.78	11.78	11.78	11.82	4.97
$d/2$	10	100	5	48.34	20.41	20.43	20.41	20.32	20.41	14.78
$d/2$	10	100	10	46.39	18.04	18.03	18.04	18.04	18.04	14.34
$d/2$	10	100	20	41.44	21.48	21.50	21.48	21.47	21.51	20.72
d	10	100	5	48.34	33.09	33.42	33.44	33.44	33.33	30.27
d	10	100	10	46.39	38.06	38.15	38.05	38.05	38.05	37.57
d	10	100	20	41.44	34.17	34.34	34.17	34.17	34.14	34.24
All	10	100	5	48.34	59.39	59.26	59.39	59.39	61.62	62.42
All	10	100	10	46.39	66.66	66.34	66.70	66.72	70.65	69.18
All	10	100	20	41.44	73.00	72.65	72.95	72.95	78.89	74.19
1	20	100	5	21.59	7.09	7.09	7.09	6.79	7.10	2.61
1	20	100	10	16.18	6.96	6.96	6.96	6.96	6.96	2.31
1	20	100	20	15.90	5.75	5.76	5.75	5.75	5.80	1.76
$d/2$	20	100	5	21.59	11.86	11.94	11.86	11.86	11.86	10.55
$d/2$	20	100	10	16.18	13.46	13.49	13.46	13.46	13.45	12.45
$d/2$	20	100	20	15.90	15.56	15.62	15.56	15.56	15.59	14.97
d	20	100	5	21.59	23.09	23.12	23.09	23.08	23.11	22.70
d	20	100	10	16.18	22.91	22.92	22.90	22.89	22.91	22.27
d	20	100	20	15.90	23.27	23.30	23.27	23.27	23.30	22.76
All	20	100	5	21.59	61.64	61.46	61.61	61.66	64.91	65.80
All	20	100	10	16.18	71.06	70.78	71.05	71.03	76.61	74.09
All	20	100	20	15.90	77.22	76.85	77.13	77.20	83.29	78.79

Table 1: Benchmarks obtained from different relaxations

From the experiments, we can see that, when using a single choice of $\mathcal{D}_\kappa$ (i.e. $|\mathcal{K}| = 1$), the choice of relaxation R_k has very little impact in the final bound. Moreover, one can see that the choice of κ matters, as choosing $\kappa = d$ has much more impact in the bound than choosing $\kappa = 1$. This makes sense, since the constraint for $\kappa = d$ considers all cumulative demands from the first until the last period, so strengthening $\mathcal{D}_\kappa$ potentially impacts previous periods implicitly as well. On the other extreme, for $\kappa = 1$, the only period involved in $\mathcal{D}_\kappa$ is the first one and so strengthening it will not affect any other future periods.

Moreover, we can observe a significant gain in bound by use of any of the relaxations when $\mathcal{K} = [d]$. How-ever, no significant advantage comes from using any of the relaxations R1, R3 or R4 instead of R2 as proposed in [12]. Thus, at least for the problem and instances in consideration, the best approach would probably be to stick to the inequalities proposed in [12]. However, one can note that using relaxation R5 allows a significant additional gap to be closed (up to 6% extra). Though significant, it is not clear if such improvement is big enough to justify a big effort to implement any cut separation using relaxation R5 more efficiently for practical purposes. However, we note that in some problems, closing 6% extra gap may be the difference to allow the solution of a

problem in reasonable time. In any case, if this extra gap closed is indeed crucial, then our results point to where to look for new cuts next.

The comparison with the strengthened star inequalities is also very interesting. When looking at the results for $|\mathcal{K}| = 1$, we see that if we just focus on $\kappa = 1$, compared to the strengthened stars, the inequalities that can be derived from the results in this paper close a lot more gap. However, when one considers $\kappa = d/2$, or $\kappa = d$, the impact of the extra effort is quite diminished. In fact, considering $\mathcal{K} = [d]$, one sees that the impact of the extra inequalities is much smaller (sometimes even worse than using just strengthened stars).

These results show the usefulness of our results in that they allow us to set the appropriate benchmarks and point to a definitive answer in terms of which direction to pursue to find useful inequalities. One comment that is worth making is that it is not contradictory that the strengthened star inequalities close more gaps than the inequalities proposed using relaxations $R_1, \ldots, R_5$ since the strengthened star inequalities are derived using integrality of the variables directly, whereas the framework proposed uses linear relaxations of the knapsack (so no integrality is directly used).

7 Conclusion

In this paper, we studied several different properties of the mixing set with a knapsack constraint and identified the key difficulty in describing its convex hull, namely that one must be able to describe the convex hull of the knapsack polytope. Since that problem by itself is hard, we were able to devise a theory that allows one to use any polyhedral relaxation of the knapsack polytope and translate that knowledge into a similar polyhedral knowledge about the desired set. We derived extended formulations for both the set of points defined by this relaxation and the cuts obtained by them. Moreover, these results generalize the results obtained in two previous papers [12, 14] dealing with the same set and we also were able to generalize a particular class of facet-defining inequalities for such set.

Finally, our computational experiments show how these theoretical results can be used to set benchmarks for identifying useful relaxations for particular chance-constrained programs.

Acknowledgement

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