

Arc connectivity and submodular flows in digraphs

Ahmad Abdi, Gérard Cornuéjols, and Giacomo Zambelli

London School of Economics and Political Science

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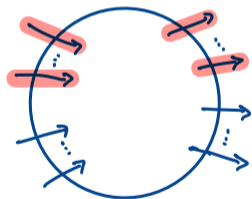
¹Paper available on the speaker's website

k -arc-connected flips

k -arc-connected flips

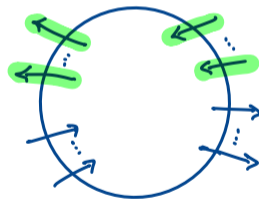
- Let $D = (V, A)$ be a digraph.
- Let $k \geq 1$ be an integer.
- A **k -arc-connected flip** is a $J \subseteq A$ such that $(D \setminus J) \cup J^{-1}$ is (strongly) k -arc-connected.

J 



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J^{-1} 

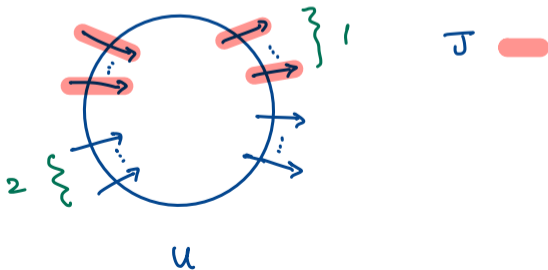


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k -arc-connected flips

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- A **k -arc-connected flip** is a $J \subseteq A$ such that $(D \setminus J) \cup J^{-1}$ is (strongly) k -arc-connected. That is,

$$\underbrace{|J \cap \delta^+(U)|}_1 + \underbrace{|\delta^-(U)| - |J \cap \delta^-(U)|}_2 \geq k \quad \forall U \subsetneq V, U \neq \emptyset.$$

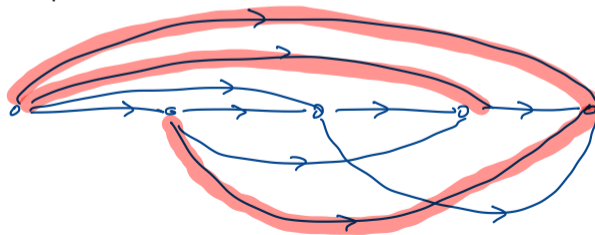


k -arc-connected flips

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- A 2-arc-connected flip:

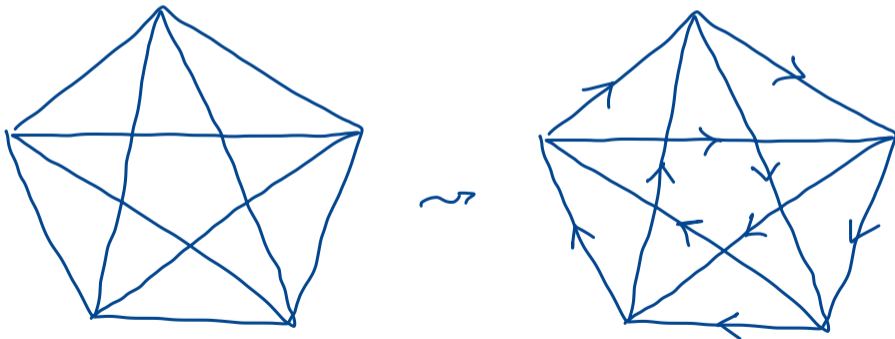


Weak Orientation Theorem

An important theorem in Graph Orientations and Submodular Optimization:

Theorem 1 (Nash-Williams '69)

If the underlying graph of D is $2k$ -edge-connected, then D has a k -arc-connected flip.



Weak Orientation Theorem

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Flash summary of the talk

- 1 We extend Theorem 1 by finding a k -arc-connected flip whose incidence vector is also a **submodular flow**.
- 2 This is made possible by finding capacitated integral solutions to the intersection of **two** submodular flow systems.

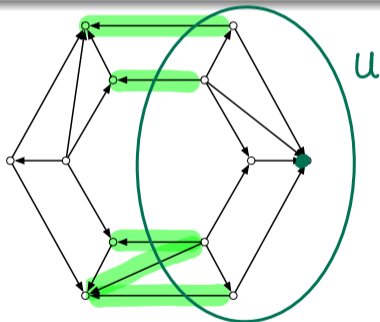
Woodall's conjecture

Dicuts

Let $D = (V, A)$ be a digraph.

Definition

A **dicut** is a subset of the form $\delta^+(U) \subseteq A$ where $U \neq \emptyset, V$ and $\delta^-(U) = \emptyset$.



Remark

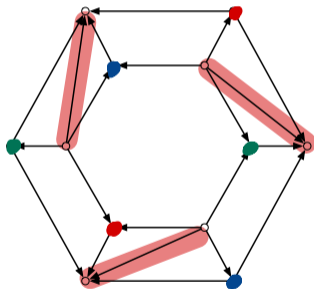
D has a dicut $\Leftrightarrow D$ is not strongly connected.

Dijoins

Let $D = (V, A)$ be a digraph.

Definition

A **dijoin** is a $J \subseteq A$ that intersects every dicut at least once, i.e. D/J is 1-arc-connected.



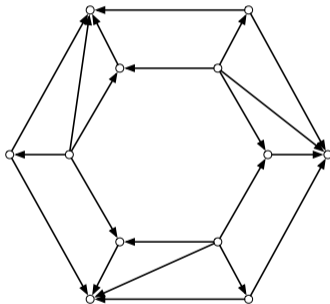
Remark

J is a dijoin $\Leftrightarrow D \cup J^{-1}$ is 1-arc-connected.

Packing dijoins

Remark

Minimum size of a **dicut** \geq maximum number of pairwise arc disjoint **dijoins**.



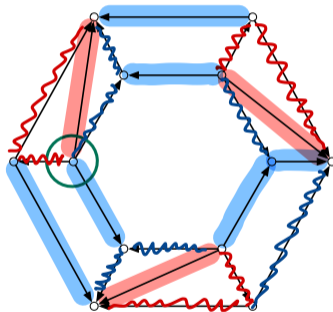
Packing dijoins

Remark

Minimum size of a **dicut** \geq maximum number of pairwise arc disjoint **dijoins**.

$$\min = 4$$

$$\max = 4$$



Packing dijoins

Remark

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Conjecture (Woodall 1978)

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Packing dijoins

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Conjecture (Woodall 1978)

Minimum size of a dicut = maximum number of pairwise arc disjoint dijoins.

- 1 **Frank and Tardos 1984**: Formulation as a common base packing problem for two matroids.
- 2 **Schrijver 1980, Feofiloff and Younger 1987**: Proved for source-sink connected digraphs.
- 3 **Lee and Wakabayashi 2001**: Proved for series-parallel digraphs.
- 4 **Mészáros 2018**: Proved for digraphs that are $(\tau - 1, 1)$ -partition-connected for τ a prime power.
- 5 **A., Cornuéjols and Zlatin 2023**: Reduced to nearly- τ -regular bipartite graphs.
- 6 Other interesting results by Cornuéjols and Guenin (2002), Shepherd and Vetta (2005), Lee and Williams (2006), Chudnovsky, Edwards, Kim, Scott, and Seymour (2016)

k -dijoins

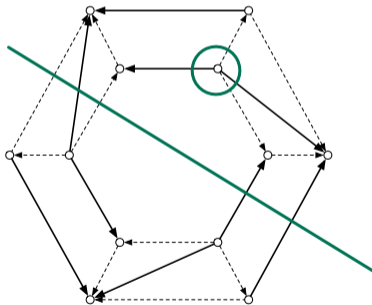
k -dijoins

Let $D = (V, A)$ be a digraph.

Definition

For an integer $k \geq 1$, a k -dijoin is an arc subset that intersects every dicut at least k times.

2-dijoin \mathcal{J} :



— \mathcal{J}

k -dijoins

Let $D = (V, A)$ be a digraph.

Definition

For an integer $k \geq 1$, a k -dijoin is an arc subset that intersects every dicut at least k times.

Proposition

The following statements hold:

- 1 The union of a disjoint k -dijoin and ℓ -dijoin is a $(k + \ell)$ -dijoin.

k -dijoins

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Proposition

The following statements hold:

- 1 The union of a disjoint k -dijoin and ℓ -dijoin is a $(k + \ell)$ -dijoin.
- 2 Every k -arc-connected flip J is a k -dijoin.

Proof of 2. Let J be a k -arc-connected flip. For every dicut $\delta^+(U)$,

$$|J \cap \delta^+(U)| = |J \cap \delta^+(U)| + |\delta^-(U)| - |J \cap \delta^-(U)| \geq k.$$

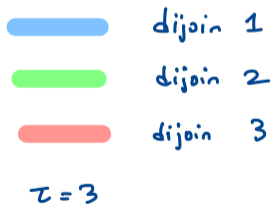
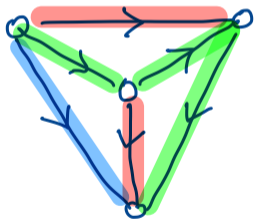
□

Decomposing into a k - and $(\tau - k)$ -dijoin

Suppose every dicut of $D = (V, A)$ has size $\geq \tau$, for some integer $\tau \geq 2$.

Conjecture (Woodall 1978)

A can be partitioned into τ dijoins.

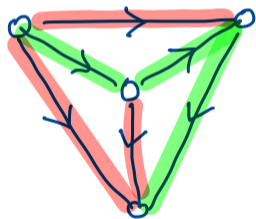


Decomposing into a k - and $(\tau - k)$ -dijoin

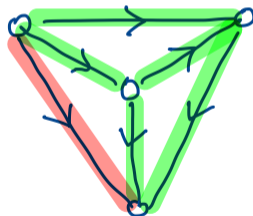
Suppose every dicut of $D = (V, A)$ has size $\geq \tau$, for some integer $\tau \geq 2$.

Conjecture ("Weak Woodall")



A can be partitioned into a k - and a $(\tau - k)$ -dijoin, for all $k \in \{1, \dots, \tau - 1\}$.



$k=1$



$k=2$

 k -dijoin
 $(\tau - k)$ -dijoin

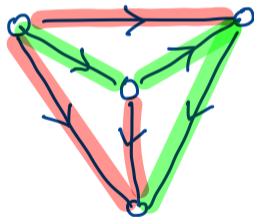
$\tau = 3$



Decomposing into a k - and $(\tau - k)$ -dijoin

Suppose every dicut of $D = (V, A)$ has size $\geq \tau$, for some integer $\tau \geq 2$.

Theorem 2 (A., Cornuéjols, Zlatin 2023)

A can be partitioned into a dijoin and a $(\tau - 1)$ -dijoin.



 1-dijoin
 $(\tau - 1)$ -dijoin

$$\tau = 3$$

Summary

Theorem 1 (Nash-Williams '69)

If the underlying graph of D is $2k$ -edge-connected, then D has a k -arc-connected flip.

Theorem 2 (A., Cornuéjols, Zlatin 2023)

If every dicut of $D = (V, A)$ has size $\geq \tau$, then A can be partitioned into a dijoin and a $(\tau - 1)$ -dijoin.

Remark

A k -arc-connected flip is a k -dijoin.

A common extension of Theorems 1 and 2

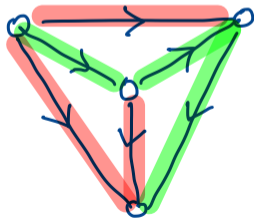
A common extension

Theorem 3 (A., Cornuéjols, Zambelli '23+)

Let $D = (V, A)$ be a digraph such that for some integers $\tau - 1 \geq k \geq 1$ we have

$$|\delta^+(U)| + \left(\frac{\tau}{k} - 1\right) |\delta^-(U)| \geq \tau \quad \forall U \subsetneq V, U \neq \emptyset.$$

Then A can be partitioned into a k -arc-connected flip and a $(\tau - k)$ -dijoin.



 1-arc-connected flip

 2-dijoin

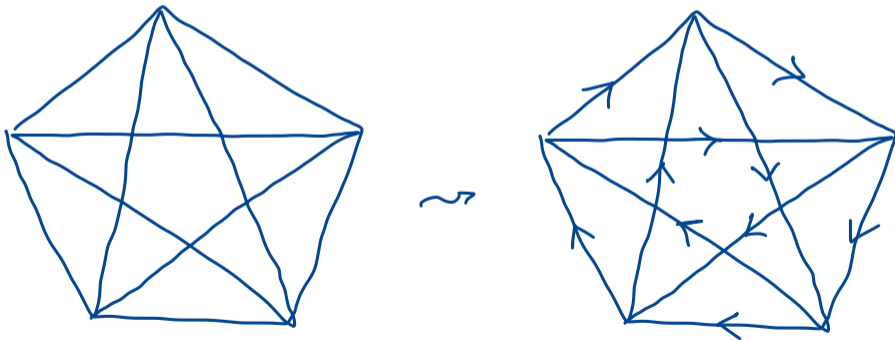
$$\tau = 3, k = 1$$

A common extension

For $\tau = 2k$ we recover:

Theorem 1 (Nash-Williams '69)

If the underlying graph of D is $2k$ -edge-connected, then D has a k -arc-connected flip.



A common extension

For $\tau = 2k$ we recover:

Theorem 1 (Nash-Williams '69)

If the underlying graph of D is $2k$ -edge-connected, then D has a k -arc-connected flip.

Why? The cut condition

$$|\delta^+(U)| + \left(\frac{\tau}{k} - 1\right) |\delta^-(U)| \geq \tau \quad \forall U \subsetneq V, U \neq \emptyset.$$

becomes

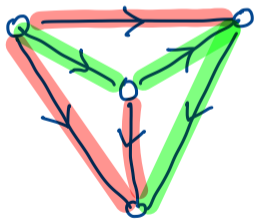
$$|\delta^+(U)| + |\delta^-(U)| \geq \tau \quad \forall U \subsetneq V, U \neq \emptyset.$$



A common extension

For $k = 1$ we recover:

Theorem 2 (A., Cornuéjols, Zlatin 2023)

If every dicut of $D = (V, A)$ has size $\geq \tau$, then A can be partitioned into a 1 -dijoin and a $(\tau - 1)$ -dijoin.



 1 -dijoin
 $(\tau - 1)$ -dijoin

$$\tau = 3$$

A common extension

For $k = 1$ we recover:

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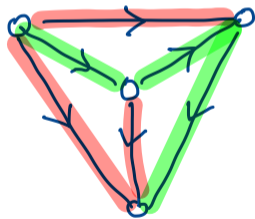
$$|\delta^+(U)| \geq \tau \quad \forall U \subsetneq V, U \neq \emptyset \text{ s.t. } \delta^-(U) = \emptyset$$

A common extension

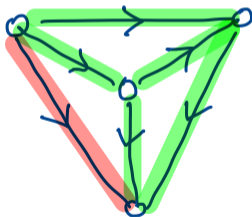
"Weak Woodall" is true for τ -edge-connected instances:

Theorem (A., Cornuéjols, Zambelli '23+)



Suppose the underlying graph of $D = (V, A)$ is τ -edge-connected. Then A can be partitioned into a k - and a $(\tau - k)$ -dijoin, for all $k \in \{1, \dots, \tau - 1\}$.



$k=1$



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 k -dijoin
 $(\tau - k)$ -dijoin

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A common extension

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Suppose the underlying graph of $D = (V, A)$ is τ -edge-connected. Then A can be partitioned into a k - and a $(\tau - k)$ -dijoin, for all $k \in \{1, \dots, \tau - 1\}$.

Proof. We may assume $\tau \geq 2k$. The cut condition holds:

$$|\delta^+(U)| + \underbrace{\left(\frac{\tau}{k} - 1\right)}_{\geq 1} |\delta^-(U)| \geq |\delta^+(U)| + |\delta^-(U)| \quad \forall U \subsetneq V, U \neq \emptyset$$
$$\geq \tau$$

\therefore \exists a partition into a k -arc-connected flip & a $(\tau - k)$ -dijoin.

□

A common extension

Theorem 3 (A., Cornuéjols, Zambelli '23+)

Let $D = (V, A)$ be a digraph such that for some integers $\tau - 1 \geq k \geq 1$ we have

$$|\delta^+(U)| + \left(\frac{\tau}{k} - 1\right) |\delta^-(U)| \geq \tau \quad \forall U \subsetneq V, U \neq \emptyset.$$

Then A can be partitioned into a k -arc-connected flip and a $(\tau - k)$ -dijoin.

Let's prove this theorem. We need two ingredients.

Ingredients

Crossing families and submodular functions

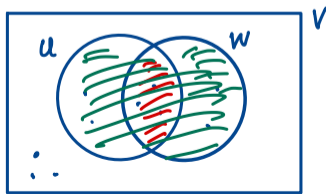
Let \mathcal{C} be a family of subsets of V , and let $f : \mathcal{C} \rightarrow \mathbb{Z}$ be a function.

Definition

\mathcal{C} is a **crossing family** if whenever $U, W \in \mathcal{C}$ and $U \cap W \neq \emptyset, U \cup W \neq V$, then $U \cap W, U \cup W \in \mathcal{C}$.

Definition

f is **crossing submodular** if whenever $U, W \in \mathcal{C}$ and $U \cap W \neq \emptyset, U \cup W \neq V$, then $f(U \cap W) + f(U \cup W) \leq f(U) + f(W)$.



Crossing families and submodular functions

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Examples

- $\{U \subsetneq V : U \neq \emptyset\}$
- $\{U \subsetneq V : U \neq \emptyset, \delta^-(U) = \emptyset\}$
- $f(U) := |\delta^+(U)|$
- $f(U) := |\delta^+(U)| - \alpha$

Ingredient 1 from Submodular Optimization

Theorem (Edmonds, Frank, Fujishige; see Schrijver 2003, Frank 2011)

Let $f_i : \mathcal{C}_i \rightarrow \mathbb{Z}$ be a crossing submodular function, for $i = 1, 2$. Then

$$\begin{aligned}x &\in \mathbb{R}^V \\ \mathbf{1}^\top x &= 0 \\ x(U) &:= \sum_{v \in U} x_v \leq f_1(U) \quad \forall U \in \mathcal{C}_1 \\ &\sum_{v \in U} x_v \leq f_2(U) \quad \forall U \in \mathcal{C}_2\end{aligned}$$

is box-totally dual integral, and therefore box-integral. In particular, **if it has a fractional solution, it has an integral solution.**

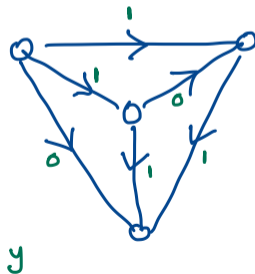
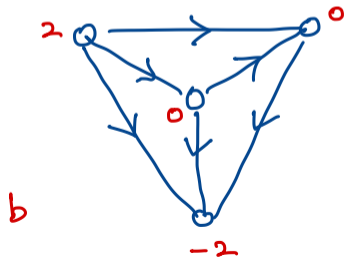
Ingredient 2 from Network Flows

Let $D = (V, A)$ be a digraph, and $b \in \mathbb{Z}^V$ s.t. $\mathbf{1}^\top b = 0$.

Definition

A **b -transshipment** is a vector $y \in \mathbb{R}^A$ s.t.

$$y(\delta^+(u)) - y(\delta^-(u)) = b_u \quad \forall u \in V.$$



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Theorem (Hoffman, Gale; see Schrijver 2003)

Take $c, d \in \mathbb{Z}^A$ with $c \leq d$ such that

$$b(U) \leq d(\delta^+(U)) - c(\delta^-(U)) \quad \forall U \neq \emptyset, V.$$

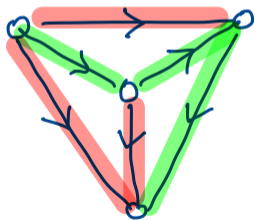
Then there exists a b -transshipment $y^* \in \mathbb{Z}^A$ such that $c \leq y^* \leq d$.

Proof of Theorem 3

Theorem 3 (A., Cornuéjols, Zambelli '23+)

A can be decomposed into a k -arc-connected flip and a $(\tau - k)$ -dijoin if for every $\emptyset \neq U \subsetneq V$,

$$|\delta^+(U)| + \left(\frac{\tau}{k} - 1\right) |\delta^-(U)| \geq \tau.$$



 1-arc-connected flip

 2-dijoin

$$\tau = 3, k = 1$$

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$$|\delta^+(U)| + \left(\frac{\tau}{k} - 1\right) |\delta^-(U)| \geq \tau.$$

Proof. We need to find a 0,1 vector y such that

- $y(\delta^+(U)) - y(\delta^-(U)) \leq |\delta^+(U)| - k$ for every $\emptyset \neq U \subsetneq V$,
- $y(\delta^+(U)) - y(\delta^-(U)) \leq |\delta^+(U)| - (\tau - k)$ for every dicut $\delta^+(U)$.

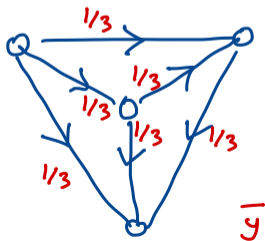
Proof continued

We know

- $|\delta^+(U)| + \left(\frac{\tau}{k} - 1\right) |\delta^-(U)| \geq \tau$ for every $\emptyset \neq U \subsetneq V$,
- Every dicut has size at least τ .

Proof. Let $\bar{y} \in \mathbb{R}^A$ assign $\frac{k}{\tau}$ to every arc. Then

- $\bar{y}(\delta^+(U)) - \bar{y}(\delta^-(U)) \leq |\delta^+(U)| - k$ for every $\emptyset \neq U \subsetneq V$,
- $\bar{y}(\delta^+(U)) - \bar{y}(\delta^-(U)) \leq |\delta^+(U)| - (\tau - k)$ for every dicut $\delta^+(U)$.



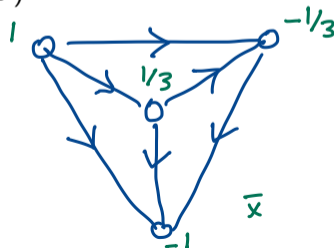
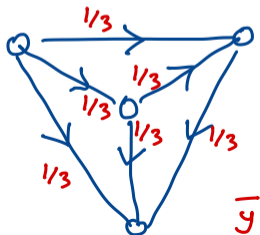
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- $\bar{y}(\delta^+(U)) - \bar{y}(\delta^-(U)) \leq |\delta^+(U)| - k$ for every $\emptyset \neq U \subsetneq V$,
- $\bar{y}(\delta^+(U)) - \bar{y}(\delta^-(U)) \leq |\delta^+(U)| - (\tau - k)$ for every dicut $\delta^+(U)$.

For each vertex v , let $\bar{x}_v := \bar{y}(\delta^+(v)) - \bar{y}(\delta^-(v))$. Then

1. $\mathbf{1}^\top \bar{x} = 0$,
2. $\bar{x}(U) \leq |\delta^+(U)| - k$ for every $\emptyset \neq U \subsetneq V$,
3. $\bar{x}(U) \leq |\delta^+(U)| - (\tau - k)$ for every dicut $\delta^+(U)$.



Proof continued

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Consequence of Ingredient 1

We can make \bar{x} integral!

Proof continued

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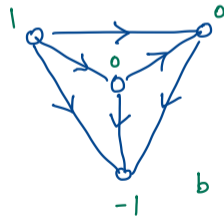
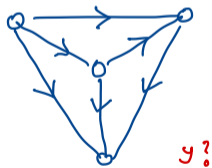
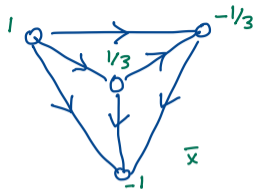
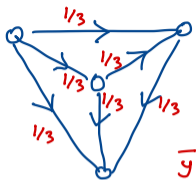
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Consequence of Ingredient 1

We can make \bar{x} integral!

But, can we make \bar{y} 0, 1?



Proof continued

There exists $b \in \mathbb{Z}^V$ such that

1. $\mathbf{1}^\top b = 0$,
2. $b(U) \leq |\delta^+(U)| - k$ for every $\emptyset \neq U \subsetneq V$,
3. $b(U) \leq |\delta^+(U)| - (\tau - k)$ for every dicut $\delta^+(U)$

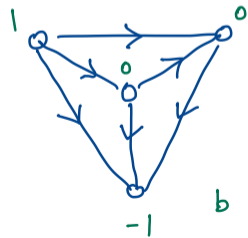
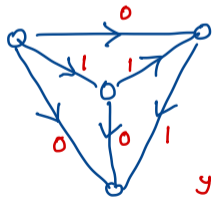
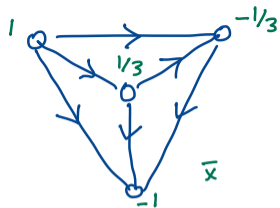
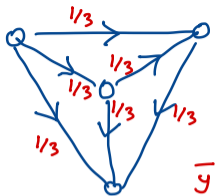
Consequence of Ingredient 2

Since $b(U) \leq |\delta^+(U)|$ for every $\emptyset \neq U \subsetneq V$, there exists a b -transshipment $y^* \in \{0, 1\}^A$, i.e. $y^*(\delta^+(v)) - y^*(\delta^-(v)) = b_v$ for every vertex v .

Then,

- $y^*(\delta^+(U)) - y^*(\delta^-(U)) \leq |\delta^+(U)| - k$ for every $\emptyset \neq U \subsetneq V$,
- $y^*(\delta^+(U)) - y^*(\delta^-(U)) \leq |\delta^+(U)| - (\tau - k)$ for every dicut $\delta^+(U)$,

as required. □



Digging a little deeper

A more powerful theorem

Theorem 4 (A., Cornuéjols, Zambelli '23+)

Let $D = (V, A)$ be a digraph, and τ, k integers with $\tau - 1 \geq k \geq 1$. Let $f : \mathcal{C} \rightarrow \mathbb{Z}$ be a crossing submodular function. Suppose the system

$$\begin{aligned}y(\delta^+(U)) - y(\delta^-(U)) &\leq |\delta^+(U)| - k && \forall U \subsetneq V, U \neq \emptyset \\y(\delta^+(U)) - y(\delta^-(U)) &\leq f(U) && \forall U \in \mathcal{C}\end{aligned}$$

is satisfied at $\bar{y} = \frac{k}{\tau} \mathbf{1}$. Then the system has a 0, 1 solution.

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For $f(U) := |\delta^+(U)| - (\tau - k)$ defined on every dicut $\delta^+(U)$, this gives

Theorem 3 (A., Cornuéjols, Zambelli '23+)

A can be decomposed into a k -arc-connected flip and a $(\tau - k)$ -dijoin if for every $\emptyset \neq U \subsetneq V$,

$$|\delta^+(U)| + \left(\frac{\tau}{k} - 1\right) |\delta^-(U)| \geq \tau.$$

Applications

For $\tau = 2k$, Theorem 4 gives

Theorem (A., Cornuéjols, Zambelli '23+)

Let $D = (V, A)$ be a digraph whose underlying graph is $2k$ -edge-connected. Let $f : \mathcal{C} \rightarrow \mathbb{Z}$ be a crossing submodular function such that

$$f(U) \geq \frac{1}{2}(|\delta^+(U)| - |\delta^-(U)|) \quad \forall U \in \mathcal{C}.$$

Then there is a k -arc-connected flip $J \subseteq A$ such that

$$f(U) \geq |J \cap \delta^+(U)| - |J \cap \delta^-(U)| \quad \forall U \in \mathcal{C}.$$

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Corollary (Nash-Williams '69)

Every $2k$ -edge-connected graph has a k -arc-connected orientation such that at every node, the out- and in-degrees differ by at most one.

Discussion

We find a 0,1 solution to

$$\begin{aligned}y(\delta^+(U)) - y(\delta^-(U)) &\leq |\delta^+(U)| - k && \forall \emptyset \neq U \subsetneq V \\y(\delta^+(U)) - y(\delta^-(U)) &\leq f(U) && \forall U \in \mathcal{C},\end{aligned}$$

the intersection of **two** 'submodular flow' systems. This is surprising...

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Folklore facts

Let $D = (V, A)$ be a digraph, and let $f_i : \mathcal{C}_i \rightarrow \mathbb{Z}$ be a crossing submodular function, for $i = 1, 2$. Consider the system

$$\begin{aligned}y &\in \mathbb{R}^A \\y(\delta^+(U)) - y(\delta^-(U)) &\leq f_1(U) \quad \forall U \in \mathcal{C}_1 \\y(\delta^+(U)) - y(\delta^-(U)) &\leq f_2(U) \quad \forall U \in \mathcal{C}_2.\end{aligned}$$

This system is **not** necessarily integral, let alone box-totally dual integral. Moreover, finding a **0,1 solution** to this system is an **NP-hard** task.

A surprising phenomenon

Intersection of two submodular flow systems

- Let $D = (V, A)$ be a digraph.
- Let $f_i : \mathcal{C}_i \rightarrow \mathbb{Z}$ be a crossing submodular function, for $i = 1, 2$.

Theorem 5 (A., Cornuéjols, Zambelli '23+)

- ① Suppose $\min_{i=1,2} f_i(U) \leq 0$ for all $U \neq \emptyset, V$ s.t. $\delta^+(U) = \delta^-(U) = \emptyset$. Then

$$y(\delta^+(U)) - y(\delta^-(U)) \leq f_1(U) \quad \forall U \in \mathcal{C}_1$$

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- ② Take $c, d \in \mathbb{Z}^A$ satisfying $c \leq d$ and

$$\min_{i=1,2} f_i(U) \leq d(\delta^+(U)) - c(\delta^-(U)) \quad \forall U \neq \emptyset, V.$$

Then every nonempty face of the feasible region contains $y^* \in \mathbb{Z}^A$ with $c \leq y^* \leq d$.

Applications

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is satisfied at $\bar{y} = \frac{k}{\tau} \mathbf{1}$. Then the system has a 0, 1 solution.

Applications

Theorem (Edmonds and Giles 1977)

Let $D = (V, A)$ be a digraph, and $f : \mathcal{C} \rightarrow \mathbb{Z}$ a crossing submodular function. The system

$$y(\delta^+(U)) - y(\delta^-(U)) \leq f(U) \quad \forall U \in \mathcal{C}$$

is box-totally dual integral, and hence box-integral.

A cute application

Theorem (A., Cornuéjols, Zambelli '23+)

Let $D = (V, A)$ be a **weakly connected** digraph. Let $f_i : \mathcal{C}_i \rightarrow \mathbb{Z}$ be a crossing submodular function, for $i = 1, 2$. Then

$$y(\delta^+(U)) - y(\delta^-(U)) \leq f_1(U) \quad \forall U \in \mathcal{C}_1$$

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This system is **not** necessarily box-integral.

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Proposition

This system is **not** necessarily box-integral.

Proposition (Goemans and Pan)

This system is **not** necessarily box-half-integral.

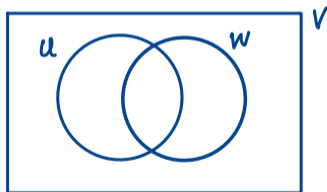
A conjecture

Tree Orientation Conjecture

Let \mathcal{C} be a family of subsets of V .

Definition

\mathcal{C} is a **lattice family** if $U \cap W, U \cup W \in \mathcal{C}$ for all $U, W \in \mathcal{C}$.



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Tree Orientation Conjecture

Let $T = (V, E)$ be a tree, and let \mathcal{C} be a lattice family over ground set V . Suppose

$$|\delta_T(U)| \geq 2 \quad \forall U \in \mathcal{C}, U \neq \emptyset, V.$$

Then there is an orientation \vec{T} of T such that

$$\delta_{\vec{T}}^+(U), \delta_{\vec{T}}^-(U) \neq \emptyset \quad \forall U \in \mathcal{C}, U \neq \emptyset, V.$$

Thanks!